Geometry Seminar Tuesday, Feb 3, 2009 Room 317 WWH at 6:00 P.M.

Small-size Epsilon-Nets for Geometric Range Spaces

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Since their introduction in 1987 by Haussler and Welzl, ϵ -nets have become one of the central concepts in computational and combinatorial geometry, and have been used in a variety of applications, such as range searching, geometric partitions, and bounds on curve-point incidences. In particular, they are strongly related to geometric set-cover and hitting-set problems.

A range space (or a hypergraph) (X,R) is a pair consisting of an underlying universe X of objects, and a certain collection R of subsets of X (also called ranges). Given a range space (X,R), and a parameter $0 < \epsilon < 1$, an ϵ -net for (X,R) is a subset N of X with the property that any range that captures at least ϵ -fraction of X contains an element of N. In other words, N is a hitting set for all the "heavy" ranges.

Of particular interest are geometric range spaces, since then they admit small-size ϵ -nets. Specifically, the Epsilon-Net Theorem of Haussler and Welzl asserts that in this case there exists an ϵ -net of size $O(1/\epsilon \log 1/\epsilon)$. One of the major questions in the theory of ϵ -nets, open since their introduction more than 20 years ago, is whether the factor $\log 1/\epsilon$ in the upper bound on their size is really necessary, especially in typical low-dimensional geometric situations. A central motivation then arises from the technique of Bronnimann and Goodrich to obtain, in polynomial time, improved approximation factors for the geometric hitting-set and set-cover problems: The existence of an ϵ -net of size $O((1/\epsilon)f(1/\epsilon))$, for some slowly-growing function $f(\ldots)$, implies an approximation factor of O(f(OPT)), where OPT is the size of the smallest such set.

In this talk I will survey some of the fundamental results concerning small-size ϵ -nets. I will then discuss range spaces of points and axis-parallel boxes in two and three dimensions, and show that they admit an ϵ -net of size $O(1/\epsilon \log \log 1/\epsilon)$, and also present several extensions to "fat" planar ranges.

Joint work with Boris Aronov (Polytechnic Institute of NYU) and Micha Sharir (Tel Aviv University).

For more information please visit the seminar website at: http://www.math.nyu.edu/seminars/geometry_seminar.html.