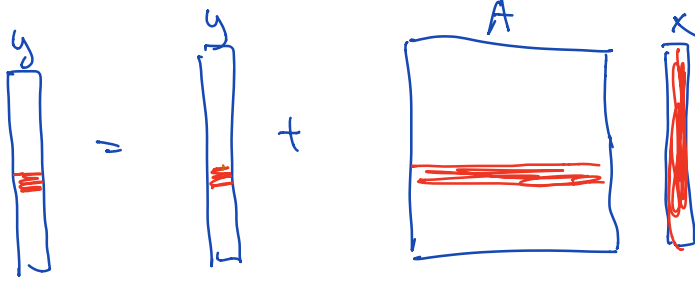


Computational intensity for matrix-vector & matrix-matrix mult.

1.) Matrix-vector multiplication

$$y = y + Ax, \quad y, x \in \mathbb{R}^n, \quad A \in \mathbb{R}^{n \times n}$$



flops:

$$\sim 2n^2 \text{ flops}$$

$$y_i = y_i + \sum_{k=1}^n a_{ik} x_k$$

$$= y_i + a_{i1}x_1 + a_{i2}x_2 + \dots + a_{in}x_n$$

memory access (slow mem access, nothing is in fast memory)
 assume we can hold a few vectors in fast memory

load y, x , load A row-by-row, write result to y

$$\underline{\underline{3n + n^2}}$$

$$\text{comp. intensity } \rho_n = f/m = \frac{2n^2}{3n + n^2} \sim \underline{\underline{2}}$$

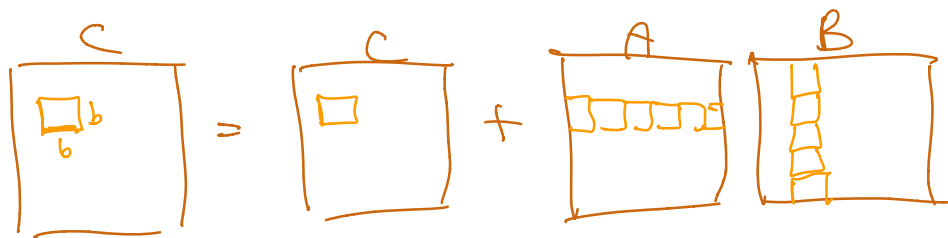
2.) Matrix-Matrix Mult

$$C = C + A * B, \quad A, B, C \in \mathbb{R}^{n \times n}$$

flops: $2n^3$

memory: same as before: $(3n + n^2)n \Rightarrow \underline{\underline{\rho_n \sim 2}}$

Tiling algorithm:



m entries, N blocks, $b = \frac{m}{N}$ blocksize

memory access: N^3 block reads of B

N^3 — w of A

$2N^2$ block reads in C
& writes

memory access: $(2N^3 + 2N^2) \cdot b^2 =$

$$(2N^3 + 2N^2) \cdot \frac{m^2}{N^2} \sim 2Nn^2 + 2n^2$$

$$q_n^{-1} = \frac{2 \frac{n^3}{b} + 2n^2}{2n^3} = 2 \frac{n^3}{b} + 2n^2$$

$$\Rightarrow q_n = b$$

\Rightarrow gives a much higher comp. intensity, much faster!