



Randomness

DS-GA 1013 / MATH-GA 2824 Optimization-based Data Analysis

http://www.cims.nyu.edu/~cfgranda/pages/OBDA_fall17/index.html

Carlos Fernandez-Granda

Gaussian random variables

Gaussian random vectors

Randomized projections

SVD of a random matrix

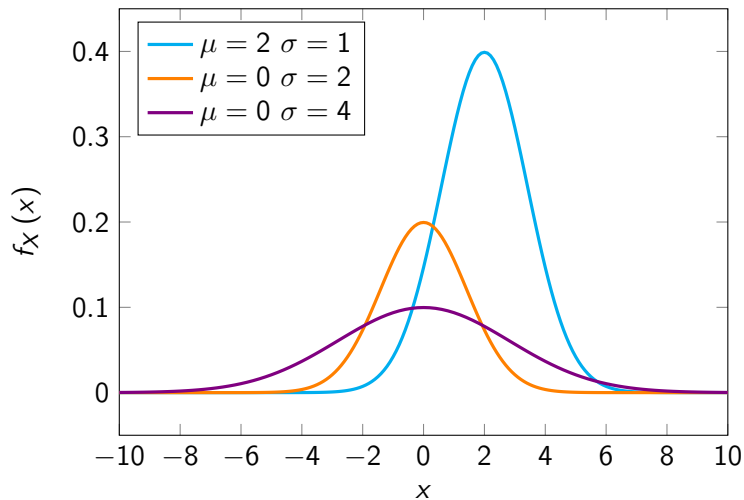
Randomized SVD

Gaussian random variables

The pdf of a Gaussian or normal random variable with mean μ and standard deviation σ is given by

$$f_X(x) = \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

Gaussian random variables



Linear transformation of Gaussian

If x is a Gaussian random variable with mean μ and standard deviation σ , then for any $a, b \in \mathbb{R}$

$$y := ax + b$$

is a Gaussian random variable with mean $a\mu + b$ and standard deviation $|a|\sigma$

Proof

Let $a > 0$ (proof for $a < 0$ is very similar), to

$$F_y(y)$$

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$$F_{\mathbf{y}}(y) = P(\mathbf{y} \leq y)$$

Proof

Let $a > 0$ (proof for $a < 0$ is very similar), to

$$\begin{aligned}F_{\mathbf{y}}(y) &= \mathbb{P}(\mathbf{y} \leq y) \\ &= \mathbb{P}(ax + b \leq y)\end{aligned}$$

Proof

Let $a > 0$ (proof for $a < 0$ is very similar), to

$$\begin{aligned}F_{\mathbf{y}}(y) &= \mathbb{P}(\mathbf{y} \leq y) \\&= \mathbb{P}(a\mathbf{x} + b \leq y) \\&= \mathbb{P}\left(\mathbf{x} \leq \frac{y - b}{a}\right)\end{aligned}$$

Proof

Let $a > 0$ (proof for $a < 0$ is very similar), to

$$\begin{aligned}F_{\mathbf{y}}(y) &= \mathbb{P}(\mathbf{y} \leq y) \\&= \mathbb{P}(ax + b \leq y) \\&= \mathbb{P}\left(\mathbf{x} \leq \frac{y - b}{a}\right) \\&= \int_{-\infty}^{\frac{y-b}{a}} \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(x-\mu)^2}{2\sigma^2}} dx\end{aligned}$$

Proof

Let $a > 0$ (proof for $a < 0$ is very similar), to

$$\begin{aligned}F_{\mathbf{y}}(y) &= \mathbb{P}(\mathbf{y} \leq y) \\&= \mathbb{P}(a\mathbf{x} + b \leq y) \\&= \mathbb{P}\left(\mathbf{x} \leq \frac{y-b}{a}\right) \\&= \int_{-\infty}^{\frac{y-b}{a}} \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(x-\mu)^2}{2\sigma^2}} dx \\&= \int_{-\infty}^y \frac{1}{\sqrt{2\pi}a\sigma} e^{-\frac{(w-a\mu-b)^2}{2a^2\sigma^2}} dw \quad \text{change of variables } w = ax + b\end{aligned}$$

Proof

Let $a > 0$ (proof for $a < 0$ is very similar), to

$$\begin{aligned}F_{\mathbf{y}}(y) &= \mathbb{P}(\mathbf{y} \leq y) \\&= \mathbb{P}(ax + b \leq y) \\&= \mathbb{P}\left(\mathbf{x} \leq \frac{y - b}{a}\right) \\&= \int_{-\infty}^{\frac{y-b}{a}} \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(x-\mu)^2}{2\sigma^2}} dx \\&= \int_{-\infty}^y \frac{1}{\sqrt{2\pi}a\sigma} e^{-\frac{(w-a\mu-b)^2}{2a^2\sigma^2}} dw \quad \text{change of variables } w = ax + b\end{aligned}$$

Differentiating with respect to y :

$$f_{\mathbf{y}}(y) = \frac{1}{\sqrt{2\pi}a\sigma} e^{-\frac{(w-a\mu-b)^2}{2a^2\sigma^2}}$$

Central limit theorem

Let x_1, x_2, x_3, \dots be a sequence of iid random variables with mean μ and bounded variance σ^2

The sequence of averages a_1, a_2, a_3, \dots is defined as

$$\mathbf{a}_i := \frac{1}{i} \sum_{j=1}^i \mathbf{x}_j$$

Central limit theorem

The sequence $\mathbf{b}_1, \mathbf{b}_2, \mathbf{b}_3, \dots$

$$\mathbf{b}_i := \sqrt{i}(\mathbf{a}_i - \mu)$$

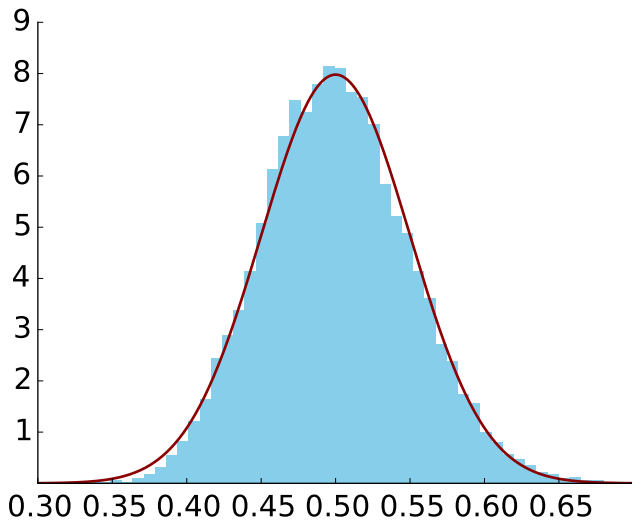
converges in distribution to a Gaussian random variable with mean 0 and variance σ^2

For any $x \in \mathbb{R}$

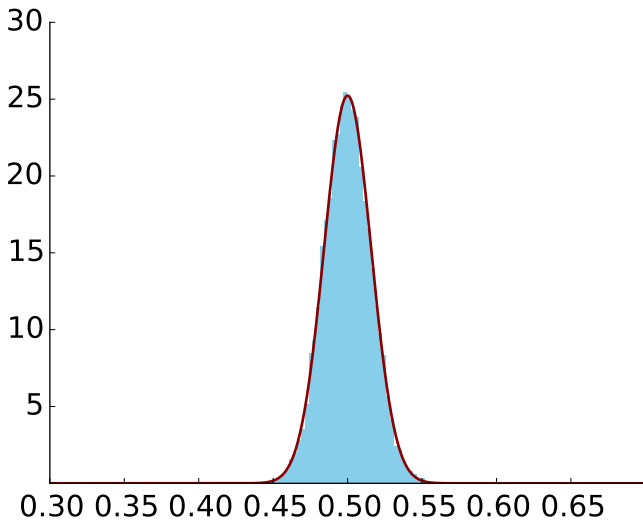
$$\lim_{i \rightarrow \infty} f_{\mathbf{b}_i}(x) := \frac{1}{\sqrt{2\pi\sigma}} e^{-\frac{x^2}{2\sigma^2}}$$

For large i the theorem suggests that the average \mathbf{a}_i is approximately **Gaussian** with mean μ and variance σ/\sqrt{n}

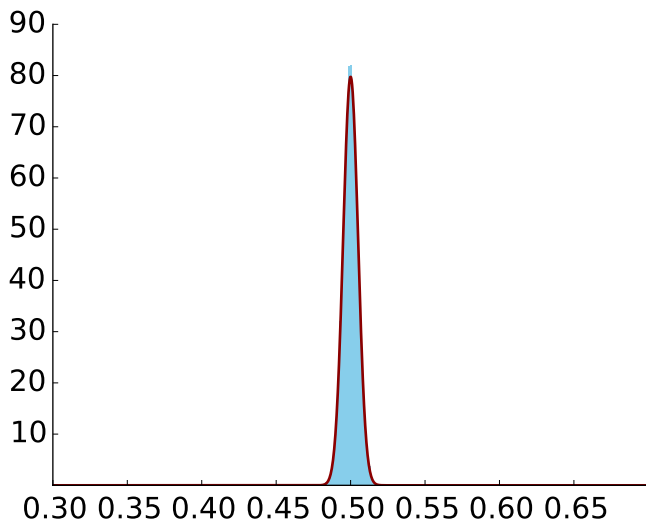
iid exponential $\lambda = 2$, $i = 10^2$



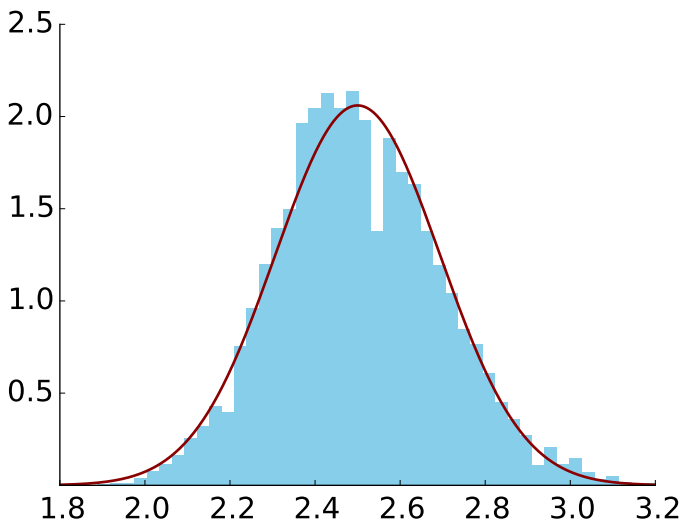
iid exponential $\lambda = 2$, $i = 10^3$



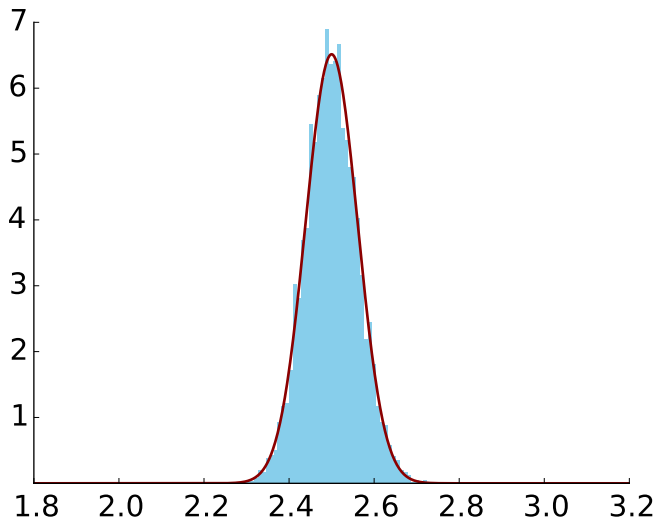
iid exponential $\lambda = 2$, $i = 10^4$



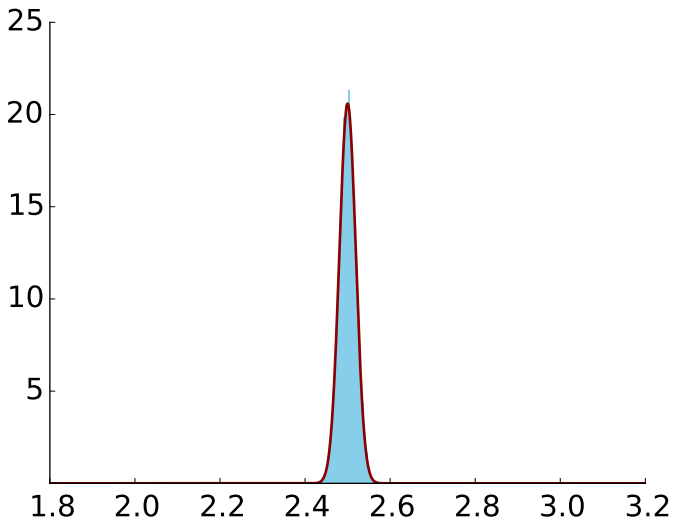
iid geometric $p = 0.4$, $i = 10^2$



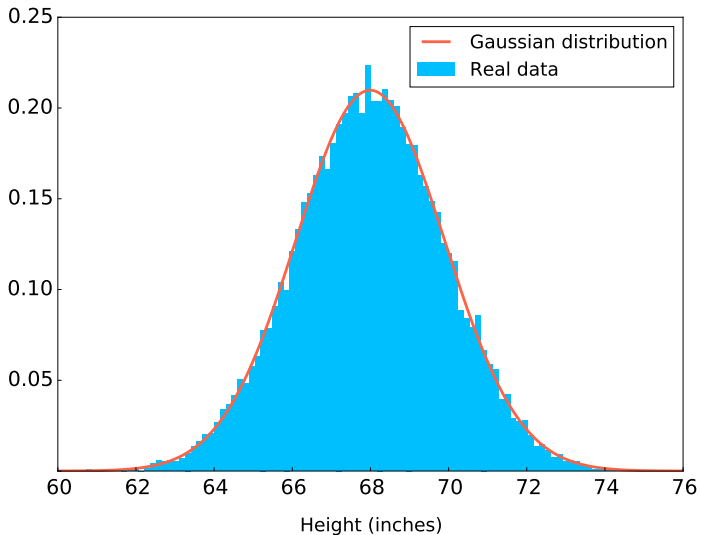
iid geometric $p = 0.4$, $i = 10^3$



iid geometric $p = 0.4$, $i = 10^4$



Histogram of heights



Gaussian random variables

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Gaussian random vector

A Gaussian random vector \vec{x} is a random vector with joint pdf

$$f_{\vec{x}}(\vec{x}) = \frac{1}{\sqrt{(2\pi)^n |\Sigma|}} \exp\left(-\frac{1}{2} (\vec{x} - \vec{\mu})^T \Sigma^{-1} (\vec{x} - \vec{\mu})\right)$$

where $\vec{\mu} \in \mathbb{R}^n$ is the mean and $\Sigma \in \mathbb{R}^{n \times n}$ the covariance matrix

Uncorrelation implies independence

If the covariance matrix is diagonal,

$$\Sigma_{\vec{x}} = \begin{bmatrix} \sigma_1^2 & 0 & \cdots & 0 \\ 0 & \sigma_2^2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & \sigma_n^2 \end{bmatrix},$$

the entries are independent

Proof

$$\Sigma_{\vec{x}}^{-1} = \begin{bmatrix} \frac{1}{\sigma_1^2} & 0 & \cdots & 0 \\ 0 & \frac{1}{\sigma_2^2} & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & \frac{1}{\sigma_n^2} \end{bmatrix}$$

$$|\Sigma| = \prod_{i=1}^n \sigma_i^2$$

Proof

$$f_{\vec{x}}(\vec{x})$$

Proof

$$f_{\vec{x}}(\vec{x}) = \frac{1}{\sqrt{(2\pi)^n |\Sigma|}} \exp\left(-\frac{1}{2}(\vec{x} - \vec{\mu})^T \Sigma^{-1}(\vec{x} - \vec{\mu})\right)$$

Proof

$$\begin{aligned} f_{\vec{x}}(\vec{x}) &= \frac{1}{\sqrt{(2\pi)^n |\Sigma|}} \exp\left(-\frac{1}{2}(\vec{x} - \vec{\mu})^T \Sigma^{-1}(\vec{x} - \vec{\mu})\right) \\ &= \prod_{i=1}^n \frac{1}{\sqrt{(2\pi)\sigma_i}} \exp\left(-\frac{(\vec{x}_i - \mu_i)^2}{2\sigma_i^2}\right) \end{aligned}$$

Proof

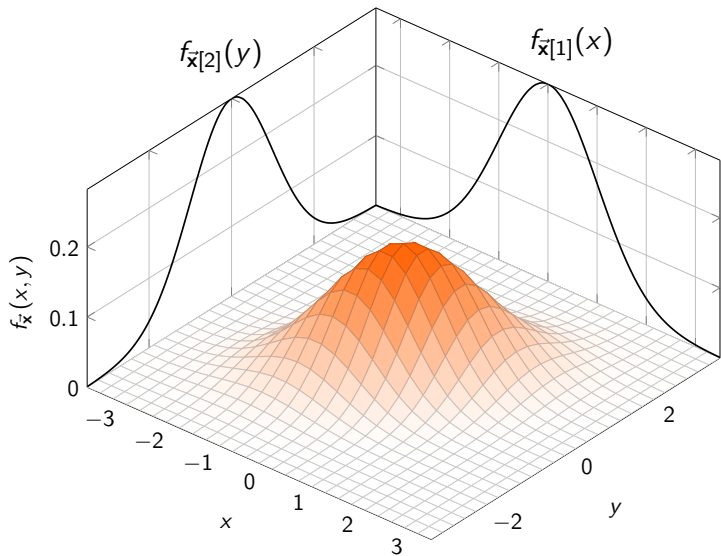
$$\begin{aligned}f_{\vec{x}}(\vec{x}) &= \frac{1}{\sqrt{(2\pi)^n |\Sigma|}} \exp\left(-\frac{1}{2}(\vec{x} - \vec{\mu})^T \Sigma^{-1}(\vec{x} - \vec{\mu})\right) \\&= \prod_{i=1}^n \frac{1}{\sqrt{(2\pi)\sigma_i}} \exp\left(-\frac{(\vec{x}_i - \mu_i)^2}{2\sigma_i^2}\right) \\&= \prod_{i=1}^n f_{\vec{x}_i}(\vec{x}_i)\end{aligned}$$

Linear transformations

Let \vec{x} be a Gaussian random vector of dimension n with mean $\vec{\mu}$ and covariance matrix Σ

For any matrix $A \in \mathbb{R}^{m \times n}$ and $\vec{b} \in \mathbb{R}^m$ $\vec{Y} = A\vec{x} + \vec{b}$ is **Gaussian** with mean $A\vec{\mu} + \vec{b}$ and covariance matrix $A\Sigma A^T$

Subvectors are also Gaussian



Direction of iid standard Gaussian vectors

If the covariance matrix of a Gaussian vector \vec{x} is I , then \vec{x} is **isotropic**

It does not favor any direction

For any orthogonal matrix U $U\vec{x}$ has the same distribution
(Gaussian with mean $U\vec{0} = \vec{0}$ and covariance matrix $UU^T = I$)

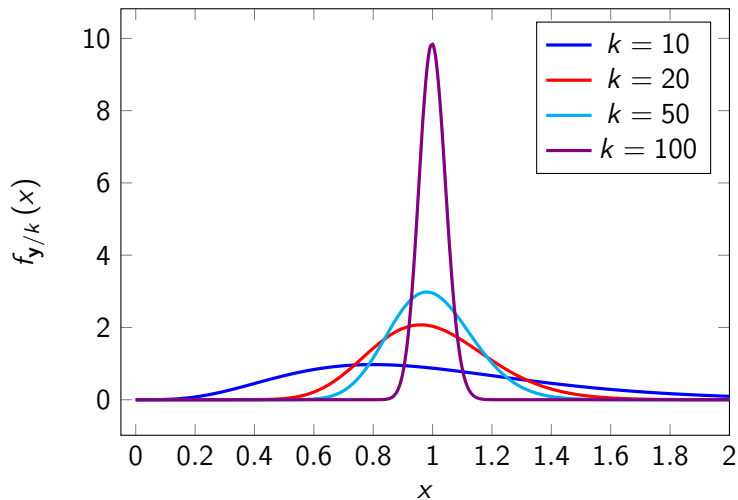
Magnitude of iid standard Gaussian vectors

In low dimensions joint pdf is mostly concentrated around the origin

High dimensions?

$\|\vec{x}\|_2^2 = \sum_{i=1}^k \vec{x}[i]^2$ is a χ^2 (chi squared) random variable with k degrees of freedom

Magnitude of iid standard Gaussian vectors



Mean

$$E\left(\|\bar{\mathbf{x}}\|_2^2\right)$$

Mean

$$\mathbb{E} \left(\|\vec{x}\|_2^2 \right) = \mathbb{E} \left(\sum_{i=1}^k \vec{x}[i]^2 \right)$$

Mean

$$\begin{aligned}\mathbb{E} \left(\|\bar{\mathbf{x}}\|_2^2 \right) &= \mathbb{E} \left(\sum_{i=1}^k \bar{\mathbf{x}}[i]^2 \right) \\ &= \sum_{i=1}^k \mathbb{E} \left(\bar{\mathbf{x}}[i]^2 \right)\end{aligned}$$

Mean

$$\begin{aligned} \mathbb{E} \left(\|\bar{\mathbf{x}}\|_2^2 \right) &= \mathbb{E} \left(\sum_{i=1}^k \bar{\mathbf{x}}[i]^2 \right) \\ &= \sum_{i=1}^k \mathbb{E} \left(\bar{\mathbf{x}}[i]^2 \right) \\ &= k \end{aligned}$$

Variance

$$\mathbb{E} \left(\left(\|\vec{x}\|_2^2 \right)^2 \right)$$

Variance

$$\mathbb{E} \left(\left(\|\vec{x}\|_2^2 \right)^2 \right) = \mathbb{E} \left(\left(\sum_{i=1}^k \vec{x}[i]^2 \right)^2 \right)$$

Variance

$$\begin{aligned} \mathbb{E} \left(\left(\|\vec{x}\|_2^2 \right)^2 \right) &= \mathbb{E} \left(\left(\sum_{i=1}^k \vec{x}[i]^2 \right)^2 \right) \\ &= \mathbb{E} \left(\sum_{i=1}^k \sum_{j=1}^k \vec{x}[i]^2 \vec{x}[j]^2 \right) \end{aligned}$$

Variance

$$\begin{aligned} \mathbb{E} \left(\left(\|\vec{x}\|_2^2 \right)^2 \right) &= \mathbb{E} \left(\left(\sum_{i=1}^k \vec{x}[i]^2 \right)^2 \right) \\ &= \mathbb{E} \left(\sum_{i=1}^k \sum_{j=1}^k \vec{x}[i]^2 \vec{x}[j]^2 \right) \\ &= \sum_{i=1}^k \sum_{j=1}^k \mathbb{E} (\vec{x}[i]^2 \vec{x}[j]^2) \end{aligned}$$

Variance

$$\begin{aligned} \mathbb{E} \left(\left(\|\bar{\mathbf{x}}\|_2^2 \right)^2 \right) &= \mathbb{E} \left(\left(\sum_{i=1}^k \bar{\mathbf{x}}[i]^2 \right)^2 \right) \\ &= \mathbb{E} \left(\sum_{i=1}^k \sum_{j=1}^k \bar{\mathbf{x}}[i]^2 \bar{\mathbf{x}}[j]^2 \right) \\ &= \sum_{i=1}^k \sum_{j=1}^k \mathbb{E} (\bar{\mathbf{x}}[i]^2 \bar{\mathbf{x}}[j]^2) \\ &= \sum_{i=1}^k \mathbb{E} (\bar{\mathbf{x}}[i]^4) + 2 \sum_{i=1}^{k-1} \sum_{j=i}^k \mathbb{E} (\bar{\mathbf{x}}[i]^2) \mathbb{E} (\bar{\mathbf{x}}[j]^2) \end{aligned}$$

Variance

$$\begin{aligned} \mathbb{E} \left(\left(\|\vec{x}\|_2^2 \right)^2 \right) &= \mathbb{E} \left(\left(\sum_{i=1}^k \vec{x}[i]^2 \right)^2 \right) \\ &= \mathbb{E} \left(\sum_{i=1}^k \sum_{j=1}^k \vec{x}[i]^2 \vec{x}[j]^2 \right) \\ &= \sum_{i=1}^k \sum_{j=1}^k \mathbb{E} (\vec{x}[i]^2 \vec{x}[j]^2) \\ &= \sum_{i=1}^k \mathbb{E} (\vec{x}[i]^4) + 2 \sum_{i=1}^{k-1} \sum_{j=i}^k \mathbb{E} (\vec{x}[i]^2) \mathbb{E} (\vec{x}[j]^2) \\ &= 3k + k(k-1) \quad \text{4th moment of standard Gaussian equals 3} \end{aligned}$$

Variance

$$\begin{aligned} \mathbb{E} \left(\left(\|\vec{x}\|_2^2 \right)^2 \right) &= \mathbb{E} \left(\left(\sum_{i=1}^k \vec{x}[i]^2 \right)^2 \right) \\ &= \mathbb{E} \left(\sum_{i=1}^k \sum_{j=1}^k \vec{x}[i]^2 \vec{x}[j]^2 \right) \\ &= \sum_{i=1}^k \sum_{j=1}^k \mathbb{E} (\vec{x}[i]^2 \vec{x}[j]^2) \\ &= \sum_{i=1}^k \mathbb{E} (\vec{x}[i]^4) + 2 \sum_{i=1}^{k-1} \sum_{j=i}^k \mathbb{E} (\vec{x}[i]^2) \mathbb{E} (\vec{x}[j]^2) \\ &= 3k + k(k-1) \quad \text{4th moment of standard Gaussian equals 3} \\ &= k(k+2) \end{aligned}$$

Variance

$$\begin{aligned}\text{Var} \left(\|\bar{\mathbf{x}}\|_2^2 \right) &= \mathbb{E} \left(\left(\|\bar{\mathbf{x}}\|_2^2 \right)^2 \right) - \mathbb{E} \left(\|\bar{\mathbf{x}}\|_2^2 \right)^2 \\ &= k(k+2) - k^2 = 2k\end{aligned}$$

Relative standard deviation around mean scales as $\sqrt{2/k}$

Non-asymptotic tail bound

Let \vec{x} be an iid standard Gaussian random vector of dimension k

For any $\epsilon > 0$

$$P\left(k(1 - \epsilon) < \|\vec{x}\|_2^2 < k(1 + \epsilon)\right) \geq 1 - \frac{2}{k\epsilon^2}$$

Markov's inequality

Let x be a **nonnegative** random variable

For any positive constant $a > 0$,

$$P(x \geq a) \leq \frac{E(x)}{a}$$

Proof

Define the indicator variable $1_{x \geq a}$

$$x - a 1_{x \geq a} \geq 0$$

Proof

Define the indicator variable $1_{\mathbf{x} \geq a}$

$$\mathbf{x} - a 1_{\mathbf{x} \geq a} \geq 0$$

$$E(\mathbf{x}) \geq a E(1_{\mathbf{x} \geq a}) = a P(\mathbf{x} \geq a)$$

Chebyshev bound

Let $\mathbf{y} := \|\vec{\mathbf{x}}\|_2^2$,

$$P(|\mathbf{y} - k| \geq k\epsilon)$$

Chebyshev bound

Let $\mathbf{y} := \|\vec{\mathbf{x}}\|_2^2$,

$$P(|\mathbf{y} - k| \geq k\epsilon) = P\left((\mathbf{y} - E(\mathbf{y}))^2 \geq k^2\epsilon^2\right)$$

Chebyshev bound

Let $\mathbf{y} := \|\vec{\mathbf{x}}\|_2^2$,

$$\begin{aligned} P(|\mathbf{y} - k| \geq k\epsilon) &= P\left((\mathbf{y} - E(\mathbf{y}))^2 \geq k^2\epsilon^2\right) \\ &\leq \frac{E\left((\mathbf{y} - E(\mathbf{y}))^2\right)}{k^2\epsilon^2} \quad \text{by Markov's inequality} \end{aligned}$$

Chebyshev bound

Let $\mathbf{y} := \|\vec{\mathbf{x}}\|_2^2$,

$$\begin{aligned} P(|\mathbf{y} - k| \geq k\epsilon) &= P\left((\mathbf{y} - E(\mathbf{y}))^2 \geq k^2\epsilon^2\right) \\ &\leq \frac{E\left((\mathbf{y} - E(\mathbf{y}))^2\right)}{k^2\epsilon^2} && \text{by Markov's inequality} \\ &= \frac{\text{Var}(\mathbf{y})}{k^2\epsilon^2} \end{aligned}$$

Chebyshev bound

Let $\mathbf{y} := \|\vec{\mathbf{x}}\|_2^2$,

$$\begin{aligned} P(|\mathbf{y} - k| \geq k\epsilon) &= P\left((\mathbf{y} - E(\mathbf{y}))^2 \geq k^2\epsilon^2\right) \\ &\leq \frac{E\left((\mathbf{y} - E(\mathbf{y}))^2\right)}{k^2\epsilon^2} && \text{by Markov's inequality} \\ &= \frac{\text{Var}(\mathbf{y})}{k^2\epsilon^2} \\ &= \frac{2}{k\epsilon^2} \end{aligned}$$

Non-asymptotic Chernoff tail bound

Let \vec{x} be an iid standard Gaussian random vector of dimension k

For any $\epsilon > 0$

$$P\left(k(1 - \epsilon) < \|\vec{x}\|_2^2 < k(1 + \epsilon)\right) \geq 1 - 2 \exp\left(-\frac{k\epsilon^2}{8}\right)$$

Proof

Let $\mathbf{y} := \|\bar{\mathbf{x}}\|_2^2$. The result is implied by

$$P(\mathbf{y} > k(1 + \epsilon)) \leq \exp\left(-\frac{k\epsilon^2}{8}\right)$$

$$P(\mathbf{y} < k(1 - \epsilon)) \leq \exp\left(-\frac{k\epsilon^2}{8}\right)$$

Proof

Fix $t > 0$

$P(\mathbf{y} > a)$

Proof

Fix $t > 0$

$$P(\mathbf{y} > a) = P(\exp(t\mathbf{y}) > \exp(at))$$

Proof

Fix $t > 0$

$$\begin{aligned} P(\mathbf{y} > a) &= P(\exp(\mathbf{t}\mathbf{y}) > \exp(\mathbf{a}\mathbf{t})) \\ &\leq \exp(-\mathbf{a}\mathbf{t}) \mathbb{E}(\exp(\mathbf{t}\mathbf{y})) \quad \text{by Markov's inequality} \end{aligned}$$

Proof

Fix $t > 0$

$$\begin{aligned} P(\mathbf{y} > a) &= P(\exp(t\mathbf{y}) > \exp(at)) \\ &\leq \exp(-at) \mathbb{E}(\exp(t\mathbf{y})) \quad \text{by Markov's inequality} \\ &\leq \exp(-at) \mathbb{E}\left(\exp\left(\sum_{i=1}^k t\mathbf{x}_i^2\right)\right) \end{aligned}$$

Proof

Fix $t > 0$

$$\begin{aligned} P(\mathbf{y} > a) &= P(\exp(t\mathbf{y}) > \exp(at)) \\ &\leq \exp(-at) \mathbb{E}(\exp(t\mathbf{y})) \quad \text{by Markov's inequality} \\ &\leq \exp(-at) \mathbb{E}\left(\exp\left(\sum_{i=1}^k t\mathbf{x}_i^2\right)\right) \\ &\leq \exp(-at) \prod_{i=1}^k \mathbb{E}(\exp(t\mathbf{x}_i^2)) \quad \text{by independence of } \mathbf{x}_1, \dots, \mathbf{x}_k \end{aligned}$$

Proof

Lemma (by direct integration)

$$E(\exp(tx^2)) = \frac{1}{\sqrt{1-2t}}$$

Equivalent to controlling higher-order moments since

$$\begin{aligned} E(\exp(tx^2)) &= E\left(\sum_{i=0}^{\infty} \frac{(tx^2)^i}{i!}\right) \\ &= \sum_{i=0}^{\infty} \frac{E(t^i (x^{2i}))}{i!}. \end{aligned}$$

Proof

Fix $t > 0$

$$\begin{aligned} P(\mathbf{y} > a) &\leq \exp(-at) \prod_{i=1}^k \mathbb{E}(\exp(tx_i^2)) \\ &= \frac{\exp(-at)}{(1-2t)^{\frac{k}{2}}} \end{aligned}$$

Proof

Setting $a := k(1 + \epsilon)$ and

$$t := \frac{1}{2} - \frac{1}{2(1 + \epsilon)},$$

we conclude

$$\begin{aligned} P(\mathbf{y} > k(1 + \epsilon)) &\leq (1 + \epsilon)^k 2 \exp\left(-\frac{k\epsilon}{2}\right) \\ &\leq \exp\left(-\frac{k\epsilon^2}{8}\right) \end{aligned}$$

Projection onto a fixed subspace

$\mathcal{P}_{S_1} \vec{z}$



$\mathcal{P}_{S_2} \vec{z}$



$$0.007 = \frac{\|\mathcal{P}_{S_1} \vec{z}\|_2}{\|\vec{x}\|_2} < \frac{\|\mathcal{P}_{S_2} \vec{z}\|_2}{\|\vec{x}\|_2} = 0.043$$

$$\frac{0.043}{0.007} = 6.14 \approx \sqrt{\frac{\dim(S_2)}{\dim(S_1)}} \quad (\text{not a coincidence})$$

Projection onto a fixed subspace

Let \mathcal{S} be a k -dimensional subspace of \mathbb{R}^n and $\vec{z} \in \mathbb{R}^n$ a vector of iid standard Gaussian noise

$\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2$ is a χ^2 random variable with k degrees of freedom

It has the same distribution as

$$y := \sum_{i=1}^k x_i^2$$

where x_1, \dots, x_k are iid standard Gaussians.

Proof

Let UU^T be a projection matrix for \mathcal{S} , where the columns of $U \in \mathbb{R}^{n \times k}$ are orthonormal:

$$\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2$$

Proof

Let UU^T be a projection matrix for \mathcal{S} , where the columns of $U \in \mathbb{R}^{n \times k}$ are orthonormal:

$$\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2 = \|UU^T \vec{z}\|_2^2$$

Proof

Let UU^T be a projection matrix for \mathcal{S} , where the columns of $U \in \mathbb{R}^{n \times k}$ are orthonormal:

$$\begin{aligned}\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2 &= \left\| UU^T \vec{z} \right\|_2^2 \\ &= \vec{z}^T UU^T UU^T \vec{z}\end{aligned}$$

Proof

Let UU^T be a projection matrix for \mathcal{S} , where the columns of $U \in \mathbb{R}^{n \times k}$ are orthonormal:

$$\begin{aligned}\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2 &= \left\| UU^T \vec{z} \right\|_2^2 \\ &= \vec{z}^T UU^T UU^T \vec{z} \\ &= \vec{z}^T UU^T \vec{z}\end{aligned}$$

Proof

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$$\begin{aligned}\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2 &= \left\| UU^T \vec{z} \right\|_2^2 \\ &= \vec{z}^T UU^T UU^T \vec{z} \\ &= \vec{z}^T UU^T \vec{z} \\ &= \vec{w}^T \vec{w}\end{aligned}$$

Proof

Let UU^T be a projection matrix for \mathcal{S} , where the columns of $U \in \mathbb{R}^{n \times k}$ are orthonormal:

$$\begin{aligned}\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2 &= \left\| UU^T \vec{z} \right\|_2^2 \\ &= \vec{z}^T UU^T UU^T \vec{z} \\ &= \vec{z}^T UU^T \vec{z} \\ &= \vec{w}^T \vec{w} \\ &= \sum_{i=1}^k \vec{w}[i]^2\end{aligned}$$

Proof

Let UU^T be a projection matrix for \mathcal{S} , where the columns of $U \in \mathbb{R}^{n \times k}$ are orthonormal:

$$\begin{aligned}\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2 &= \left\| UU^T \vec{z} \right\|_2^2 \\ &= \vec{z}^T UU^T UU^T \vec{z} \\ &= \vec{z}^T UU^T \vec{z} \\ &= \vec{w}^T \vec{w} \\ &= \sum_{i=1}^k \vec{w}[i]^2\end{aligned}$$

$\vec{w} := U^T \vec{z}$ is Gaussian with mean zero and covariance matrix

$$\Sigma_{\vec{w}} = U^T \Sigma_{\vec{z}} U$$

Proof

Let UU^T be a projection matrix for \mathcal{S} , where the columns of $U \in \mathbb{R}^{n \times k}$ are orthonormal:

$$\begin{aligned}\|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2^2 &= \left\| UU^T \vec{z} \right\|_2^2 \\ &= \vec{z}^T UU^T UU^T \vec{z} \\ &= \vec{z}^T UU^T \vec{z} \\ &= \vec{w}^T \vec{w} \\ &= \sum_{i=1}^k \vec{w}[i]^2\end{aligned}$$

$\vec{w} := U^T \vec{z}$ is Gaussian with mean zero and covariance matrix

$$\begin{aligned}\Sigma_{\vec{w}} &= U^T \Sigma_{\vec{z}} U \\ &= U^T U = I\end{aligned}$$

Non-asymptotic Chernoff tail bound

Let \vec{x} be an iid standard Gaussian random vector of dimension k

For any $\epsilon > 0$

$$P\left(k(1 - \epsilon) < \|\vec{x}\|_2^2 < k(1 + \epsilon)\right) \geq 1 - 2 \exp\left(-\frac{k\epsilon^2}{8}\right)$$

Projection onto a fixed subspace

Let \mathcal{S} be a k -dimensional subspace of \mathbb{R}^n and $\vec{z} \in \mathbb{R}^n$ a vector of iid standard Gaussian noise

For any $\epsilon > 0$

$$P(k(1 - \epsilon) < \|\mathcal{P}_{\mathcal{S}} \vec{z}\|_2 < k(1 + \epsilon)) \geq 1 - 2 \exp\left(-\frac{k\epsilon^2}{8}\right)$$

Gaussian random variables

Gaussian random vectors

Randomized projections

SVD of a random matrix

Randomized SVD

Dimensionality reduction

- ▶ PCA preserves the most *energy* (ℓ_2 norm)
- ▶ **Problem 1:** Computationally expensive
- ▶ **Problem 2:** Depends on *all* of the data
- ▶ (Possible) **Solution:** Just project randomly!
- ▶ For a data set $\vec{x}_1, \vec{x}_2, \dots \in \mathbb{R}^m$ compute $\mathbf{A}\vec{x}_1, \mathbf{A}\vec{x}_2, \dots \in \mathbb{R}^m$ where $\mathbf{A} \in \mathbb{R}^{k \times n}$ ($k < n$) has iid standard Gaussian entries

Fixed vector

Let \mathbf{A} be a $a \times b$ matrix with iid standard Gaussian entries

If $\vec{v} \in \mathbb{R}^b$ is a deterministic vector with unit ℓ_2 norm, then $\mathbf{A}\vec{v}$ is an **a -dimensional iid standard Gaussian** vector

Proof:

Fixed vector

Let \mathbf{A} be a $a \times b$ matrix with iid standard Gaussian entries

If $\vec{v} \in \mathbb{R}^b$ is a deterministic vector with unit ℓ_2 norm, then $\mathbf{A}\vec{v}$ is an **a -dimensional iid standard Gaussian** vector

Proof:

$(\mathbf{A}\vec{v})[i]$, $1 \leq i \leq a$ is Gaussian with mean zero and variance

$$\begin{aligned}\text{Var}(\mathbf{A}_{i,:}^T \vec{v}) &= \vec{v}^T \Sigma_{\mathbf{A}_{i,:}} \vec{v} \\ &= \vec{v}^T \mathbf{I} \vec{v} \\ &= \|\vec{v}\|_2^2 = 1\end{aligned}$$

Non-asymptotic Chernoff tail bound

Let \vec{x} be an iid standard Gaussian random vector of dimension k

For any $\epsilon > 0$

$$P\left(k(1 - \epsilon) < \|\vec{x}\|_2^2 < k(1 + \epsilon)\right) \geq 1 - 2 \exp\left(-\frac{k\epsilon^2}{8}\right)$$

Fixed vector

Let \mathbf{A} be a $a \times b$ matrix with iid standard Gaussian entries

For any $\vec{v} \in \mathbb{R}^p$ with unit norm and any $\epsilon \in (0, 1)$

$$\sqrt{a(1 - \epsilon)} \leq \|\mathbf{A}\vec{v}\|_2 \leq \sqrt{a(1 + \epsilon)}$$

with probability at least $1 - 2 \exp(-a\epsilon^2/8)$

Johnson-Lindenstrauss lemma

Let \mathbf{A} be a $k \times n$ matrix with iid standard Gaussian entries

Let $\vec{x}_1, \dots, \vec{x}_p \in \mathbb{R}^n$ be any fixed set of p deterministic vectors

For any pair \vec{x}_i, \vec{x}_j and any $\epsilon \in (0, 1)$

$$(1 - \epsilon) \|\vec{x}_i - \vec{x}_j\|_2^2 \leq \left\| \frac{1}{\sqrt{k}} \mathbf{A} \vec{x}_i - \frac{1}{\sqrt{k}} \mathbf{A} \vec{x}_j \right\|_2^2 \leq (1 + \epsilon) \|\vec{x}_i - \vec{x}_j\|_2^2$$

with probability at least $\frac{1}{p}$ as long as

$$k \geq \frac{16 \log(p)}{\epsilon^2}$$

Proof

Aim: Control action of \mathbf{A} the normalized differences

$$\vec{v}_{ij} := \frac{\vec{x}_i - \vec{x}_j}{\|\vec{x}_i - \vec{x}_j\|_2}$$

Our event of interest is the intersection of the events

$$\mathcal{E}_{ij} = \left\{ k(1 - \epsilon) < \|\mathbf{A}\vec{v}_{ij}\|_2^2 < k(1 + \epsilon) \right\} \quad 1 \leq i < p, i < j \leq p$$

Fixed vector

Let \mathbf{A} be a $a \times b$ matrix with iid standard Gaussian entries

For any $\vec{v} \in \mathbb{R}^b$ with unit norm and any $\epsilon \in (0, 1)$

$$\sqrt{a(1-\epsilon)} \leq \|\mathbf{A}\vec{v}\|_2 \leq \sqrt{a(1+\epsilon)}$$

with probability at least $1 - 2 \exp(-a\epsilon^2/8)$

This implies

$$\mathbb{P}(\mathcal{E}_{ij}^c) \leq \frac{2}{p^2} \quad \text{if } k \geq \frac{16 \log(p)}{\epsilon^2}$$

Union bound

For any events S_1, S_2, \dots, S_n in a probability space

$$P(\cup_i S_i) \leq \sum_{i=1}^n P(S_i).$$

Proof

Number of events \mathcal{E}_{ij} equals $\binom{p}{2} = p(p-1)/2$

By the union bound

$$P\left(\bigcap_{i,j} \mathcal{E}_{ij}\right)$$

Proof

Number of events \mathcal{E}_{ij} equals $\binom{p}{2} = p(p-1)/2$

By the union bound

$$P\left(\bigcap_{i,j} \mathcal{E}_{ij}\right) = 1 - P\left(\bigcup_{i,j} \mathcal{E}_{ij}^c\right)$$

Proof

Number of events \mathcal{E}_{ij} equals $\binom{p}{2} = p(p-1)/2$

By the union bound

$$\begin{aligned} P\left(\bigcap_{i,j} \mathcal{E}_{ij}\right) &= 1 - P\left(\bigcup_{i,j} \mathcal{E}_{ij}^c\right) \\ &\geq 1 - \sum_{i,j} P(\mathcal{E}_{ij}^c) \end{aligned}$$

Proof

Number of events \mathcal{E}_{ij} equals $\binom{p}{2} = p(p-1)/2$

By the union bound

$$\begin{aligned} P\left(\bigcap_{i,j} \mathcal{E}_{ij}\right) &= 1 - P\left(\bigcup_{i,j} \mathcal{E}_{ij}^c\right) \\ &\geq 1 - \sum_{i,j} P(\mathcal{E}_{ij}^c) \\ &\geq 1 - \frac{p(p-1)}{2} \frac{2}{p^2} \end{aligned}$$

Proof

Number of events \mathcal{E}_{ij} equals $\binom{p}{2} = p(p-1)/2$

By the union bound

$$\begin{aligned} \mathbb{P} \left(\bigcap_{i,j} \mathcal{E}_{ij} \right) &= 1 - \mathbb{P} \left(\bigcup_{i,j} \mathcal{E}_{ij}^c \right) \\ &\geq 1 - \sum_{i,j} \mathbb{P} (\mathcal{E}_{ij}^c) \\ &\geq 1 - \frac{p(p-1)}{2} \frac{2}{p^2} \\ &\geq \frac{1}{p} \end{aligned}$$

Dimensionality reduction for visualization

Motivation: Visualize high-dimensional features projected onto 2D or 3D

Example:

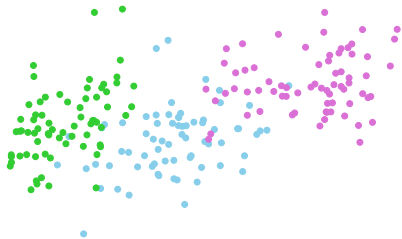
Seeds from three different varieties of wheat: Kama, Rosa and Canadian

Features:

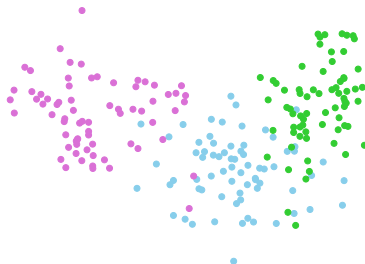
- ▶ Area
- ▶ Perimeter
- ▶ Compactness
- ▶ Length of kernel
- ▶ Width of kernel
- ▶ Asymmetry coefficient
- ▶ Length of kernel groove

Dimensionality reduction for visualization

Randomized projection



PCA



Nearest neighbors in random subspace

Nearest neighbors classification (Algorithm 4.2 in Lecture Notes 1) computes n distances in \mathbb{R}^m for each new example

Cost: $\mathcal{O}(nmp)$ for p examples

Idea: Use a $k \times m$ iid standard Gaussian matrix to project onto k -dimensional space beforehand

Cost:

- ▶ kmn operations to project training set
- ▶ kmp operations to project test set
- ▶ knp to perform nearest-neighbor classification

Much faster!

Face recognition

Training set: 360 64×64 images from 40 different subjects (9 each)

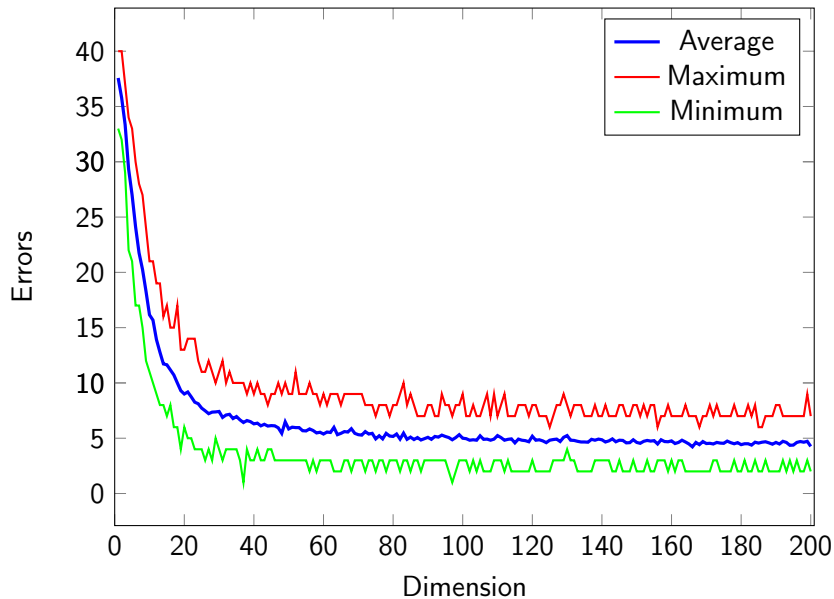
Test set: 1 new image from each subject

We model each image as a vector in \mathbb{R}^{4096} ($m = 4096$)

To classify we:

1. Project onto random a k -dimensional subspace
2. Apply nearest-neighbor classification using the ℓ_2 -norm distance in \mathbb{R}^k

Performance



Nearest neighbor in \mathbb{R}^{50}

Test image



Projection



Closest projection



Corresponding image



Gaussian random variables

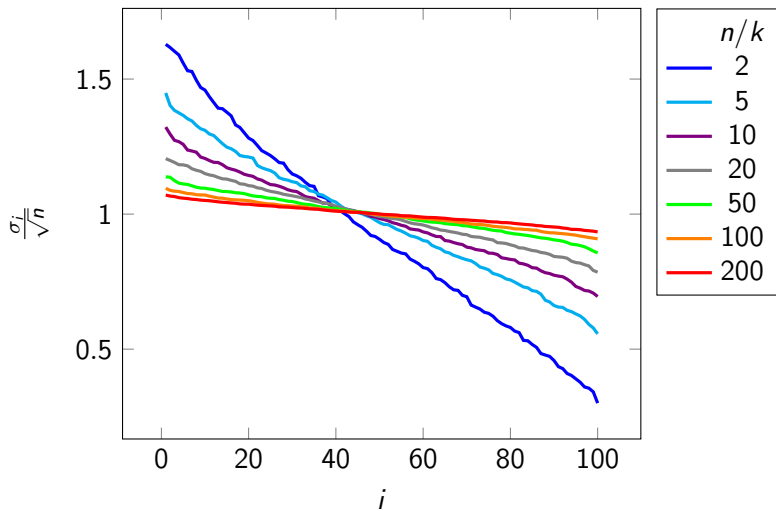
Gaussian random vectors

Randomized projections

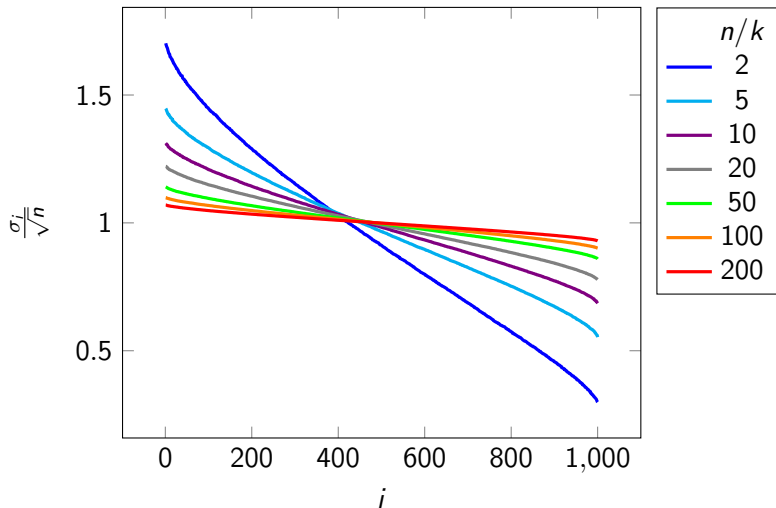
SVD of a random matrix

Randomized SVD

Singular values of $n \times k$ matrix, $k = 100$



Singular values of $n \times k$ matrix, $k = 1000$



Singular values of a Gaussian matrix

Intuitively as n grows

$$\mathbf{A} \approx U (\sqrt{n} I) V^T = \sqrt{n} UV^T,$$

iid Gaussian vectors in high dimensions are *almost* orthogonal

Singular values of a Gaussian matrix

Let \mathbf{A} be a $n \times k$ matrix with iid standard Gaussian entries such that $n > k$

For any fixed $\epsilon > 0$, the singular values of \mathbf{A} satisfy

$$\sqrt{n(1-\epsilon)} \leq \sigma_k \leq \sigma_1 \leq \sqrt{n(1+\epsilon)}$$

with probability at least $1 - 1/k$ as long as

$$n > \frac{64k}{\epsilon^2} \log \frac{12}{\epsilon}$$

Proof

Recall that

$$\sigma_1 = \max_{\{\|\vec{x}\|_2=1 \mid \vec{x} \in \mathbb{R}^k\}} \|\mathbf{A}\vec{x}\|_2$$

$$\sigma_k = \min_{\{\|\vec{x}\|_2=1 \mid \vec{x} \in \mathbb{R}^k\}} \|\mathbf{A}\vec{x}\|_2$$

so the bounds are equivalent to

$$n(1 - \epsilon) < \|\mathbf{A}\vec{v}\|_2^2 < n(1 + \epsilon)$$

Proof

Idea: Use union bound over all unit-norm vectors

Problem: They are infinite!

Solution: Use union bound on a finite set, then show that this is enough

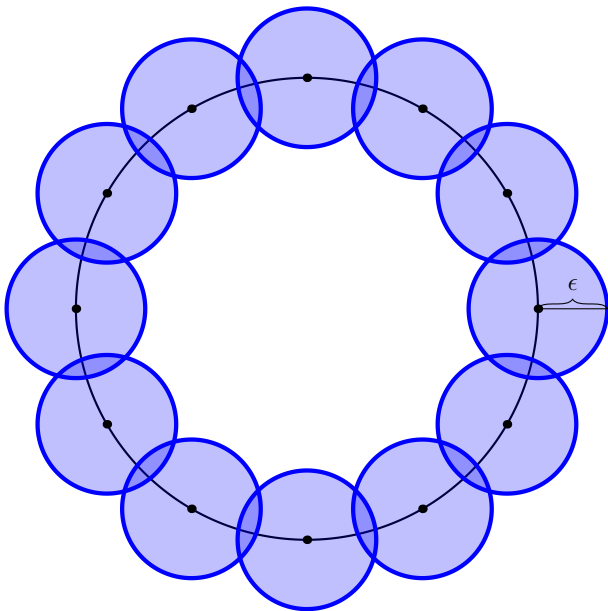
ϵ -net

An ϵ -net of a set $\mathcal{X} \subseteq \mathbb{R}^k$ is a subset $\mathcal{N}_\epsilon \subseteq \mathcal{X}$ such that for every vector $\vec{x} \in \mathcal{X}$ there exists $\vec{y} \in \mathcal{N}_\epsilon$ for which

$$\|\vec{x} - \vec{y}\|_2 \leq \epsilon.$$

The **covering number** $\mathcal{N}(\mathcal{X}, \epsilon)$ of a set \mathcal{X} at scale ϵ is the minimal cardinality of an ϵ -net of \mathcal{X}

ϵ -net



Covering number of a sphere

The covering number of the n -dimensional sphere \mathcal{S}^{k-1} at scale ϵ satisfies

$$\mathcal{N}(\mathcal{S}^{k-1}, \epsilon) \leq \left(\frac{2 + \epsilon}{\epsilon}\right)^k \leq \left(\frac{3}{\epsilon}\right)^k$$

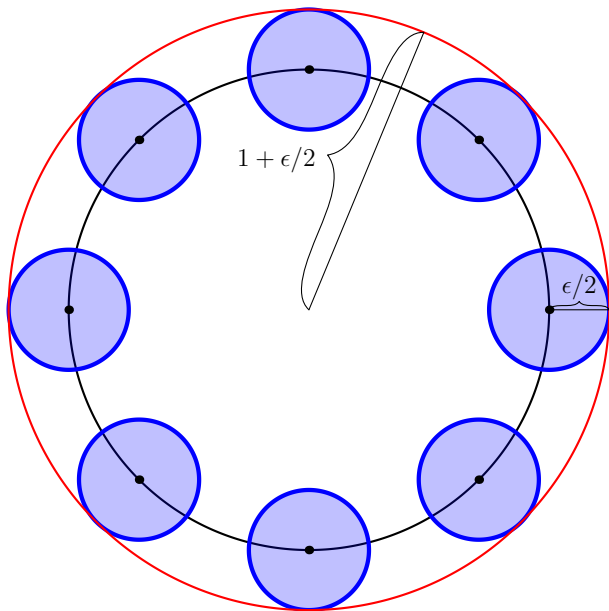
Covering number of a sphere

- ▶ Initialize \mathcal{N}_ϵ to the empty set
- ▶ Choose a point $\vec{x} \in \mathcal{S}^{k-1}$ such that

$$\|\vec{x} - \vec{y}\|_2 > \epsilon \quad \text{for any } \vec{y} \in \mathcal{N}_\epsilon$$

- ▶ Add \vec{x} to \mathcal{N}_ϵ until there are no points in \mathcal{S}^{k-1} that are ϵ away from any point in \mathcal{N}_ϵ

Covering number of a sphere



Covering number of a sphere

$$\text{Vol} \left(\mathcal{B}_{1+\epsilon/2}^k \left(\vec{0} \right) \right) \geq \text{Vol} \left(\bigcup_{\vec{x} \in \mathcal{N}_\epsilon} \mathcal{B}_{\epsilon/2}^k \left(\vec{x} \right) \right)$$

Covering number of a sphere

$$\begin{aligned}\text{Vol} \left(\mathcal{B}_{1+\epsilon/2}^k \left(\vec{0} \right) \right) &\geq \text{Vol} \left(\bigcup_{\vec{x} \in \mathcal{N}_\epsilon} \mathcal{B}_{\epsilon/2}^k \left(\vec{x} \right) \right) \\ &= |\mathcal{N}_\epsilon| \text{Vol} \left(\mathcal{B}_{\epsilon/2}^k \left(\vec{0} \right) \right)\end{aligned}$$

Covering number of a sphere

$$\begin{aligned}\text{Vol} \left(\mathcal{B}_{1+\epsilon/2}^k \left(\vec{0} \right) \right) &\geq \text{Vol} \left(\cup_{\vec{x} \in \mathcal{N}_\epsilon} \mathcal{B}_{\epsilon/2}^k \left(\vec{x} \right) \right) \\ &= |\mathcal{N}_\epsilon| \text{Vol} \left(\mathcal{B}_{\epsilon/2}^k \left(\vec{0} \right) \right)\end{aligned}$$

By multivariable calculus

$$\text{Vol} \left(\mathcal{B}_r^k \left(\vec{0} \right) \right) = r^k \text{Vol} \left(\mathcal{B}_1^k \left(\vec{0} \right) \right)$$

Covering number of a sphere

$$\begin{aligned}\text{Vol} \left(\mathcal{B}_{1+\epsilon/2}^k \left(\vec{0} \right) \right) &\geq \text{Vol} \left(\bigcup_{\vec{x} \in \mathcal{N}_\epsilon} \mathcal{B}_{\epsilon/2}^k \left(\vec{x} \right) \right) \\ &= |\mathcal{N}_\epsilon| \text{Vol} \left(\mathcal{B}_{\epsilon/2}^k \left(\vec{0} \right) \right)\end{aligned}$$

$$\text{Vol} \left(\mathcal{B}_r^k \left(\vec{0} \right) \right) = r^k \text{Vol} \left(\mathcal{B}_1^k \left(\vec{0} \right) \right)$$

so we conclude

$$(1 + \epsilon/2)^k \geq |\mathcal{N}_\epsilon| (\epsilon/2)^k$$

Proof

1. We prove the bounds

$$n(1 - \epsilon_2) < \|\mathbf{A}\vec{v}\|_2^2 < n(1 + \epsilon_2)$$

where $\epsilon_2 := \epsilon/2$ on an $\epsilon_1 := \epsilon/4$ net of the sphere

2. We show that by the triangle inequality, this implies that the bounds hold on all the sphere

Fixed vector

Let \mathbf{A} be a $a \times b$ matrix with iid standard Gaussian entries

For any $\vec{v} \in \mathbb{R}^b$ with unit norm and any $\epsilon \in (0, 1)$

$$\sqrt{a(1 - \epsilon)} \leq \|\mathbf{A}\vec{v}\|_2 \leq \sqrt{a(1 + \epsilon)}$$

with probability at least $1 - 2 \exp(-a\epsilon^2/8)$

Bound on the ϵ_1 -net

We define the event

$$\mathcal{E}_{\vec{v}, \epsilon_2} := \left\{ n(1 - \epsilon_2) \|\vec{v}\|_2^2 \leq \|\mathbf{A}\vec{v}\|_2^2 \leq n(1 + \epsilon_2) \|\vec{v}\|_2^2 \right\}$$

$$\mathbb{P} \left(\bigcup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \right)$$

Bound on the ϵ_1 -net

We define the event

$$\mathcal{E}_{\vec{v}, \epsilon_2} := \left\{ n(1 - \epsilon_2) \|\vec{v}\|_2^2 \leq \|\mathbf{A}\vec{v}\|_2^2 \leq n(1 + \epsilon_2) \|\vec{v}\|_2^2 \right\}$$

$$\mathbb{P} \left(\bigcup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \right) \leq \sum_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathbb{P} \left(\mathcal{E}_{\vec{v}, \epsilon_2}^c \right)$$

Bound on the ϵ_1 -net

We define the event

$$\mathcal{E}_{\vec{v}, \epsilon_2} := \left\{ n(1 - \epsilon_2) \|\vec{v}\|_2^2 \leq \|\mathbf{A}\vec{v}\|_2^2 \leq n(1 + \epsilon_2) \|\vec{v}\|_2^2 \right\}$$

$$\begin{aligned} \mathbb{P} \left(\bigcup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \right) &\leq \sum_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathbb{P} \left(\mathcal{E}_{\vec{v}, \epsilon_2}^c \right) \\ &\leq |\mathcal{N}_{\epsilon_1}| \mathbb{P} \left(\mathcal{E}_{\vec{v}, \epsilon_2}^c \right) \end{aligned}$$

Bound on the ϵ_1 -net

We define the event

$$\mathcal{E}_{\vec{v}, \epsilon_2} := \left\{ n(1 - \epsilon_2) \|\vec{v}\|_2^2 \leq \|\mathbf{A}\vec{v}\|_2^2 \leq n(1 + \epsilon_2) \|\vec{v}\|_2^2 \right\}$$

$$\begin{aligned} \mathbb{P} \left(\bigcup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \right) &\leq \sum_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathbb{P} \left(\mathcal{E}_{\vec{v}, \epsilon_2}^c \right) \\ &\leq |\mathcal{N}_{\epsilon_1}| \mathbb{P} \left(\mathcal{E}_{\vec{v}, \epsilon_2}^c \right) \\ &\leq 2 \left(\frac{12}{\epsilon} \right)^k \exp \left(-\frac{n\epsilon^2}{32} \right) \end{aligned}$$

Bound on the ϵ_1 -net

We define the event

$$\mathcal{E}_{\vec{v}, \epsilon_2} := \left\{ n(1 - \epsilon_2) \|\vec{v}\|_2^2 \leq \|\mathbf{A}\vec{v}\|_2^2 \leq n(1 + \epsilon_2) \|\vec{v}\|_2^2 \right\}$$

$$\begin{aligned} \mathbb{P} \left(\bigcup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \right) &\leq \sum_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathbb{P} \left(\mathcal{E}_{\vec{v}, \epsilon_2}^c \right) \\ &\leq |\mathcal{N}_{\epsilon_1}| \mathbb{P} \left(\mathcal{E}_{\vec{v}, \epsilon_2}^c \right) \\ &\leq 2 \left(\frac{12}{\epsilon} \right)^k \exp \left(-\frac{n\epsilon^2}{32} \right) \\ &\leq \frac{1}{k} \quad \text{if } n > \frac{64k}{\epsilon^2} \log \frac{12}{\epsilon} \end{aligned}$$

Upper bound on the sphere

Let $\vec{x} \in \mathcal{S}^{k-1}$

There exists $\vec{v} \in \mathcal{N}(\mathcal{X}, \epsilon_1)$ such that $\|\vec{x} - \vec{v}\|_2 \leq \epsilon/4$

$\|\mathbf{A}\vec{x}\|_2$

Upper bound on the sphere

Let $\vec{x} \in \mathcal{S}^{k-1}$

There exists $\vec{v} \in \mathcal{N}(\mathcal{X}, \epsilon_1)$ such that $\|\vec{x} - \vec{v}\|_2 \leq \epsilon/4$

$$\|\mathbf{A}\vec{x}\|_2 \leq \|\mathbf{A}\vec{v}\|_2 + \|\mathbf{A}(\vec{x} - \vec{v})\|_2$$

Upper bound on the sphere

Let $\vec{x} \in \mathcal{S}^{k-1}$

There exists $\vec{v} \in \mathcal{N}(\mathcal{X}, \epsilon_1)$ such that $\|\vec{x} - \vec{v}\|_2 \leq \epsilon/4$

$$\begin{aligned}\|\mathbf{A}\vec{x}\|_2 &\leq \|\mathbf{A}\vec{v}\|_2 + \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \\ &\leq \sqrt{n} \left(1 + \frac{\epsilon}{2}\right) + \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \quad \text{assuming } \cup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \text{ holds}\end{aligned}$$

Upper bound on the sphere

Let $\vec{x} \in \mathcal{S}^{k-1}$

There exists $\vec{v} \in \mathcal{N}(\mathcal{X}, \epsilon_1)$ such that $\|\vec{x} - \vec{v}\|_2 \leq \epsilon/4$

$$\begin{aligned}\|\mathbf{A}\vec{x}\|_2 &\leq \|\mathbf{A}\vec{v}\|_2 + \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \\ &\leq \sqrt{n} \left(1 + \frac{\epsilon}{2}\right) + \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \quad \text{assuming } \cup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \text{ holds} \\ &\leq \sqrt{n} \left(1 + \frac{\epsilon}{2}\right) + \sigma_1 \|\vec{x} - \vec{v}\|_2\end{aligned}$$

Upper bound on the sphere

Let $\vec{x} \in \mathcal{S}^{k-1}$

There exists $\vec{v} \in \mathcal{N}(\mathcal{X}, \epsilon_1)$ such that $\|\vec{x} - \vec{v}\|_2 \leq \epsilon/4$

$$\begin{aligned}\|\mathbf{A}\vec{x}\|_2 &\leq \|\mathbf{A}\vec{v}\|_2 + \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \\ &\leq \sqrt{n} \left(1 + \frac{\epsilon}{2}\right) + \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \quad \text{assuming } \cup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \text{ holds} \\ &\leq \sqrt{n} \left(1 + \frac{\epsilon}{2}\right) + \sigma_1 \|\vec{x} - \vec{v}\|_2 \\ &\leq \sqrt{n} \left(1 + \frac{\epsilon}{2}\right) + \frac{\sigma_1 \epsilon}{4}\end{aligned}$$

Upper bound on the sphere

$$\sigma_1 \leq \sqrt{n} \left(1 + \frac{\epsilon}{2}\right) + \frac{\sigma_1 \epsilon}{4}$$

$$\begin{aligned}\sigma_1 &\leq \sqrt{n} \left(\frac{1 + \epsilon/2}{1 - \epsilon/4}\right) \\ &= \sqrt{n} \left(1 + \epsilon - \frac{\epsilon(1 - \epsilon)}{4 - \epsilon}\right) \\ &\leq \sqrt{n}(1 + \epsilon)\end{aligned}$$

Lower bound on the sphere

$$\|\mathbf{A}\vec{x}\|_2$$

Lower bound on the sphere

$$\|\mathbf{A}\vec{x}\|_2 \geq \|\mathbf{A}\vec{v}\|_2 - \|\mathbf{A}(\vec{x} - \vec{v})\|_2$$

Lower bound on the sphere

$$\begin{aligned}\|\mathbf{A}\vec{x}\|_2 &\geq \|\mathbf{A}\vec{v}\|_2 - \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \|\mathbf{A}(\vec{x} - \vec{v})\|_2\end{aligned}\quad \text{assuming } \cup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c \text{ holds}$$

Lower bound on the sphere

$$\begin{aligned}\|\mathbf{A}\vec{x}\|_2 &\geq \|\mathbf{A}\vec{v}\|_2 - \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \|A(\vec{x} - \vec{v})\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \sigma_1 \|\vec{x} - \vec{v}\|_2\end{aligned}$$

assuming $\cup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c$ holds

Lower bound on the sphere

$$\begin{aligned}\|\mathbf{A}\vec{x}\|_2 &\geq \|\mathbf{A}\vec{v}\|_2 - \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \|A(\vec{x} - \vec{v})\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \sigma_1 \|\vec{x} - \vec{v}\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \frac{\epsilon}{4} \sqrt{n} (1 + \epsilon)\end{aligned}$$

assuming $\cup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c$ holds

Lower bound on the sphere

$$\begin{aligned}\|\mathbf{A}\vec{x}\|_2 &\geq \|\mathbf{A}\vec{v}\|_2 - \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \|\mathbf{A}(\vec{x} - \vec{v})\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \sigma_1 \|\vec{x} - \vec{v}\|_2 \\ &\geq \sqrt{n} \left(1 - \frac{\epsilon}{2}\right) - \frac{\epsilon}{4} \sqrt{n} (1 + \epsilon) \\ &= \sqrt{n} (1 - \epsilon)\end{aligned}$$

assuming $\cup_{\vec{v} \in \mathcal{N}_{\epsilon_1}} \mathcal{E}_{\vec{v}, \epsilon_2}^c$ holds

Gaussian random variables

Gaussian random vectors

Randomized projections

SVD of a random matrix

Randomized SVD

Fast SVD

For a matrix $M \in \mathbb{R}^{m \times n}$ which is approximately rank k :

1. Choose a small oversampling parameter p (usually 5 or slightly larger).
2. Find a matrix $\tilde{U} \in \mathbb{R}^{m \times (k+p)}$ with $k+p$ orthonormal columns that approximately span the column space of M
3. Compute $W \in \mathbb{R}^{(k+p) \times n}$ defined by $W := \tilde{U}^T M$
4. Compute the SVD of $W = U_W S_W V_W^T$
5. Output $U := (\tilde{U} U_W)_{:,1:k}$, $S := (S_W)_{1:k,1:k}$ and $V := (V_W)_{:,1:k}$ as the SVD of M

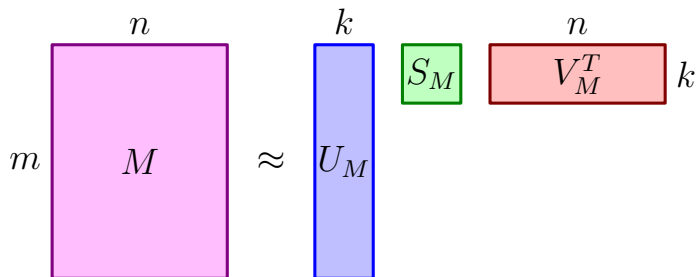
Fast SVD

For a matrix $M \in \mathbb{R}^{m \times n}$ which is approximately rank k :

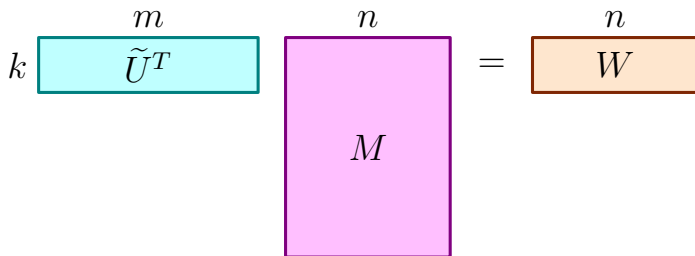
1. Choose a small oversampling parameter p (usually 5 or slightly larger).
2. Find a matrix $\tilde{U} \in \mathbb{R}^{m \times (k+p)}$ with $k+p$ orthonormal columns that approximately span the column space of M
3. Compute $W \in \mathbb{R}^{(k+p) \times n}$ defined by $W := \tilde{U}^T M \mathcal{O}(kmn)$
4. Compute the SVD of $W = U_W S_W V_W^T \mathcal{O}(k^2n)$
5. Output $U := (\tilde{U} U_W)_{:,1:k}$, $S := (S_W)_{1:k,1:k}$ and $V := (V_W)_{:,1:k}$ as the SVD of M

Complexity of regular SVD is $\mathcal{O}(mn \min\{m, n\})$

Fast SVD



Fast SVD



The diagram illustrates a matrix multiplication in the context of Fast SVD. It consists of three main components: a cyan rectangle on the left, a pink rectangle in the middle, and an orange rectangle on the right, all connected by an equals sign. The cyan rectangle is labeled with k on its left side and m above it, and contains the symbol \tilde{U}^T . The pink rectangle is labeled with n above it and contains the symbol M . The orange rectangle is labeled with n above it and contains the symbol W . The entire expression is $k \times m$ matrix \tilde{U}^T multiplied by n matrix M , resulting in an n matrix W .

$$k \times m \text{ matrix } \tilde{U}^T \times n \text{ matrix } M = n \text{ matrix } W$$

Fast SVD

The method works if (1) M is rank k and (2) \tilde{U} spans the column space

M

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$$M = \tilde{U}\tilde{U}^T M$$

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$$\begin{aligned} M &= \tilde{U}\tilde{U}^T M \\ &= \tilde{U}W \end{aligned}$$

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$$\begin{aligned}M &= \tilde{U}\tilde{U}^T M \\ &= \tilde{U}W \\ &= \tilde{U}U_W S_W V_W^T\end{aligned}$$

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where $U := \tilde{U}U_W$ is an $m \times k$ matrix with orthonormal columns

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$$U^T U = U_W^T \tilde{U}^T \tilde{U} U_W$$

Fast SVD

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where $U := \tilde{U}U_W$ is an $m \times k$ matrix with orthonormal columns

$$\begin{aligned}U^T U &= U_W^T \tilde{U}^T \tilde{U} U_W \\ &= U_W^T U_W = I\end{aligned}$$

Power iterations

For approximately low-rank matrices performance depends on **gap** between σ_k and σ_{k+1}

The gap can be increased by power iterations

This method is only used when computing \tilde{U}

The input is

$$\tilde{M} := (MM^T)^q M$$

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The input is

$$\begin{aligned}\tilde{M} &:= (MM^T)^q M \\ &= (U_M S_M^2 U_M^T)^q U_M S_M V_M^T \\ &= U_M S_M^2 U_M^T U_M S_M^2 U_M^T \cdots U_M S_M^2 U_M^T U_M V_M^T \\ &= U_M S_M^{2q+1} V_M^T\end{aligned}$$

Problem

How do we estimate the column space of a low-rank matrix?

- ▶ Project onto random subspace with slightly larger dimension
- ▶ Select random columns

Randomized column-space approximation

For a matrix $M \in \mathbb{R}^{m \times n}$ which is approximately rank k :

1. Create an $n \times (k + p)$ iid standard Gaussian matrix \mathbf{A} , where p is a small integer (e.g. 5)
2. Compute the $m \times (k + p)$ matrix $\mathbf{B} = M\mathbf{A}$
3. Orthonormalize the columns of \mathbf{B} and output them as a matrix $\tilde{\mathbf{U}} \in \mathbb{R}^{m \times (k+p)}$.
4. Apply power iterations if necessary.

Randomized column-space approximation

$$\begin{aligned}\mathbf{B} &= \mathbf{M}\mathbf{A} \\ &= \mathbf{U}_M \mathbf{S}_M \mathbf{V}_M^T \mathbf{A} \\ &= \mathbf{U}_M \mathbf{S}_M \mathbf{C}\end{aligned}$$

Randomized column-space approximation

$$\begin{aligned}\mathbf{B} &= \mathbf{M}\mathbf{A} \\ &= \mathbf{U}_M \mathbf{S}_M \mathbf{V}_M^T \mathbf{A} \\ &= \mathbf{U}_M \mathbf{S}_M \mathbf{C}\end{aligned}$$

- ▶ If M is low rank \mathbf{C} is a $k \times (k + p)$ iid standard Gaussian matrix

Randomized column-space approximation

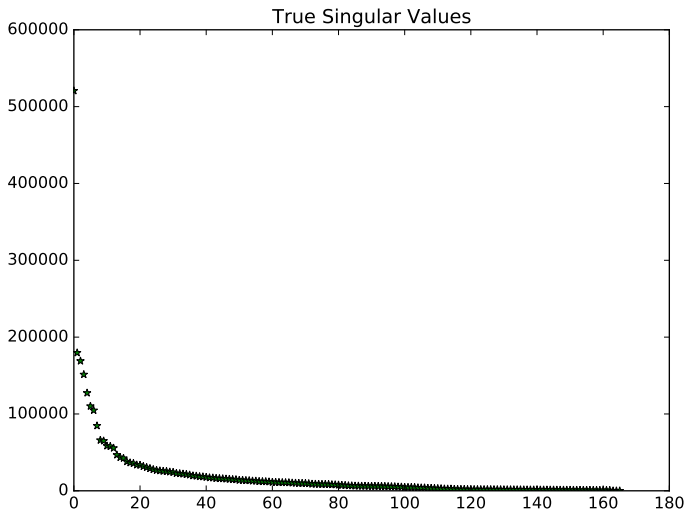
$$\begin{aligned}\mathbf{B} &= \mathbf{M}\mathbf{A} \\ &= \mathbf{U}_M \mathbf{S}_M \mathbf{V}_M^T \mathbf{A} \\ &= \mathbf{U}_M \mathbf{S}_M \mathbf{C}\end{aligned}$$

- ▶ If M is low rank \mathbf{C} is a $k \times (k + p)$ iid standard Gaussian matrix
- ▶ Otherwise, \mathbf{C} is a $\min\{m, n\} \times (k + p)$ iid standard Gaussian matrix

Randomized SVD of a video

- ▶ Video with $160 \times 1080 \times 1920$ frames
- ▶ We interpret each frame as a vector in $\mathbb{R}^{20,736,000}$
- ▶ Matrix formed by these vectors is approximately low rank
- ▶ Regular SVD takes 12 seconds (281.1 seconds if we take 691 frames)
- ▶ Fast SVD with randomized-column-space estimate takes 5.8 seconds (10.4 seconds for 691 frames) to obtain a rank-10 approximation ($q = 2, p = 7$)

True singular values



Left singular vector approximation

True



Estimated



Random column selection

For a matrix $M \in \mathbb{R}^{m \times n}$ which is approximately rank k :

1. Select a random subset of column indices $\mathcal{I} := \{\mathbf{i}_1, \mathbf{i}_2, \dots, \mathbf{i}_{k'}\}$ with $k' \geq k$
2. Orthonormalize the submatrix corresponding to \mathcal{I} :

$$M_{\mathcal{I}} := [M_{:, \mathbf{i}_1} \quad M_{:, \mathbf{i}_2} \quad \cdots \quad M_{:, \mathbf{i}_{k'}}]$$

and output them as a matrix $\tilde{\mathbf{U}} \in \mathbb{R}^{m \times k'}$

Random column selection

(Possible) **Problem**: If right singular vectors are sparse, this will not work

$$\mathbf{M}_{\mathcal{I}} = \mathbf{U}_M \mathbf{S}_M (\mathbf{V}_M)_{\mathcal{I}}$$

Example

$$M := \begin{bmatrix} -3 & 2 & 2 & 2 \\ 3 & 2 & 2 & 2 \\ -3 & 2 & 2 & 2 \\ 3 & 2 & 2 & 2 \end{bmatrix}$$

Example

$$M = U_M S_M V_M^T = \begin{bmatrix} 0.5 & -0.5 \\ 0.5 & 0.5 \\ 0.5 & -0.5 \\ 0.5 & 0.5 \end{bmatrix} \begin{bmatrix} 6.9282 & 0 \\ 0 & 6 \end{bmatrix} \begin{bmatrix} 0 & 0.577 & 0.577 & 0.577 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

Example, $\mathcal{I} = \{2, 3\}$

$$M_{\mathcal{I}} = \begin{bmatrix} 2 & 2 \\ 2 & 2 \\ 2 & 2 \\ 2 & 2 \end{bmatrix} = \begin{bmatrix} 0.5 \\ 0.5 \\ 0.5 \\ 0.5 \end{bmatrix} 6.2982 [0.577 \quad 0.577] .$$

Randomized SVD of a video

- ▶ Video with 160 1080×1920 frames
- ▶ We interpret each frame as a vector in $\mathbb{R}^{20,736,000}$
- ▶ Matrix formed by these vectors is approximately low rank
- ▶ Regular SVD takes 12 seconds (281.1 seconds if we take 691 frames)
- ▶ Fast SVD with random-column-selection estimate takes 5.2 seconds to obtain a rank-10 approximation ($k' = 17$)

Left singular vector approximation

True



Estimated



Singular value approximation

