

Theorem: $A \in \mathbb{R}^{n \times n}$ regular, $b \in \mathbb{R}^n$, $Ax = b$, $A(x + \delta x) = b + \delta b$

$$\Rightarrow \frac{\|\delta x\|}{\|x\|} \leq \kappa(A) \frac{\|\delta b\|}{\|b\|}$$

Proof: $b = Ax$, $\delta x = A^{-1}(b + \delta b) - x = A^{-1} \delta b$

$$\|b\| = \|Ax\| \leq \|A\| \|x\|, \quad \|\delta x\| \leq \|A^{-1}\| \|\delta b\|$$

multiply \rightarrow

$$\|b\| \|\delta x\| \leq \underbrace{\|A\| \|A^{-1}\|}_{\kappa(A)} \|x\| \|\delta b\|$$

$$\Rightarrow \frac{\|\delta x\|}{\|x\|} \leq \kappa(A) \frac{\|\delta b\|}{\|b\|} \quad \square.$$

Example: $A = \begin{pmatrix} 1 & \frac{1}{2} & \frac{1}{3} \\ \frac{1}{2} & \frac{1}{3} & \frac{1}{4} \\ \frac{1}{3} & \frac{1}{4} & \frac{1}{5} \end{pmatrix} \in \mathbb{R}^{3 \times 3}$, invertible,
 $\kappa_2(A) \approx 524$

We solve $Ax = b$ with $b = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$ and b are measured with potential error of 10^{-4} , i.e.

$$b + \delta b = \begin{pmatrix} 1 \\ \epsilon \\ 0 \end{pmatrix}, \quad |\epsilon| \leq 10^{-4}$$

Then x satisfies $\|x\|_2 \approx 48$, $x = \begin{pmatrix} 9 \\ -36 \\ 30 \end{pmatrix}$ error of 0.01%

$$\frac{\|\delta x\|_2}{48} \leq 524 \cdot \frac{10^{-4}}{1}$$

$$\rightarrow \|\delta x\|_2 \leq 2.5 \rightarrow \text{error in } x \text{ of } \approx 5\%$$

\rightarrow relative error potentially increases by a factor of ≈ 500 , i.e. by the condition number.

Least squares problems

Consider an overdetermined system

$$\begin{pmatrix} 3 & 1 \\ 1 & 1 \\ 4 & 2 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 2 \end{pmatrix}, \quad A \in \mathbb{R}^{m \times n}, \quad b \in \mathbb{R}^m \\ m \geq n$$

This system does, in general, not have a solution, but we can try to find $x \in \mathbb{R}^2$ such that $Ax - b \approx 0$. This can be formulated as least squares problem

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_2,$$

which is equivalent with

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_2^2 \leftarrow \text{"least squares"}$$

Since for any vector $y \in \mathbb{R}^m$: $\|y\|_2^2 = y^T y$, we can write

$$\|Ax - b\|_2^2 = (Ax - b)^T (Ax - b) = x^T A^T A x - \underbrace{(x^T A^T b + b^T A x)}_{\text{quadratic form in } x} + b^T b$$

Minimum is characterized by

$$A^T A x = A^T b \quad \text{normal equations}$$

$$\boxed{A^T} \boxed{A} \boxed{x} = \boxed{A^T} \boxed{b}$$

One should avoid computation of $A^T A$ as this can be expensive if m, n are large, and $A^T A$ can be poorly conditioned.

For instance: $A = \begin{pmatrix} \epsilon & 0 \\ 0 & 1 \end{pmatrix}, B = A^T A$

$$k_2(A) = \|A\|_2 \|A^{-1}\|_2 = 1 \cdot \epsilon^{-1} = \frac{1}{\epsilon}$$

$$k_2(B) = \|A^T A\|_2 \|(A^T A)^{-1}\|_2 = \frac{1}{\epsilon^2}$$

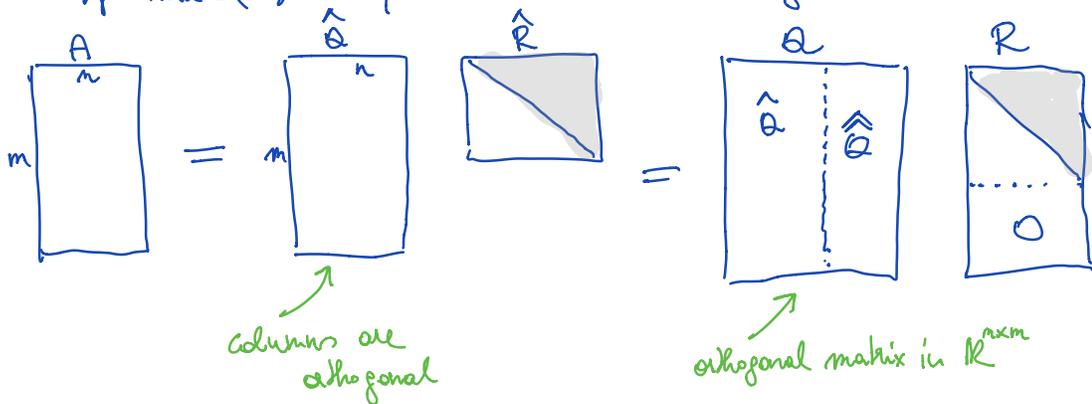
ϵ is very small, $k_2(A) = \frac{1}{\epsilon} \ll k_2(A^T A) = \frac{1}{\epsilon^2}$

Theorem: Let $A \in \mathbb{R}^{m \times n}, m \geq n$. Then

$$A = \hat{Q} \hat{R}$$

where \hat{R} is upper triangular matrix, $\hat{R} \in \mathbb{R}^{n \times n}$, $\hat{Q} \in \mathbb{R}^{m \times n}$ with $\hat{Q}^T \hat{Q} = I_{n \times n}$.

If $\text{rank}(A) = n$, then \hat{R} is non-singular.



Theorem: QR-factorization for least squares solve

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_2^2. \text{ The solution } x \text{ can be}$$

computed as solution of $\hat{R}x = \hat{Q}^T b$.

Proof: $\|Ax-b\|_2^2 = (Ax-b)^T (Ax-b)$ requires one backward subst.
 $= (Ax-b)^T \underbrace{Q Q^T}_I (Ax-b)$ $R=Q^T A$ $A=QR=\hat{Q}\hat{R}$

$$= \|Q^T (Ax-b)\|_2^2 = \|Rx - Q^T b\|_2^2$$

$$= \left\| \begin{pmatrix} \hat{R}x \\ 0 \end{pmatrix} - \begin{pmatrix} \hat{Q}^T b \\ b_2 \end{pmatrix} \right\|_2^2 = \|\hat{R}x - \hat{Q}^T b\|_2^2 +$$

$$+ \|b_2\|_2^2 \geq \|b_2\|_2^2, \quad \boxed{b_2 = \hat{Q}^T b}$$

→ The quantity $\|Ax-b\|_2^2$ can be made smallest possible by making $\|\hat{R}x - \hat{Q}^T b\|_2^2 = 0$, i.e. $\hat{R}x = \hat{Q}^T b$.

Thus, by solving $\hat{R}x = \hat{Q}^T b$, we find the solution to the least squares problem. \square .