

Power method for finding largest eigenvalue of  $A \in \mathbb{R}^{n \times n}$

Start  $x_0 \in \mathbb{R}^n$ , Compute

$$x_{k+1} = Ax_k \quad k=0, 1, 2, \dots$$

Theorem:  $\lambda_1$  is a simple eigenvalue of the symmetric matrix  $A \in \mathbb{R}^{n \times n}$

$$|\lambda_1| > |\lambda_2| \geq |\lambda_3| \dots \geq |\lambda_n|$$

Let  $x_0 \in \mathbb{R}^n$  be a vector that is not orthogonal to the eigenspace of  $\lambda_1$ . Then the sequence  $y_k := \frac{x_k}{\|x_k\|}$  with  $x_k$  defined in (\*) converges towards a normalized eigenvector of  $A$  for  $\lambda_1$ .

Proof:  $\eta_1, \dots, \eta_n$  orthonormal basis of  $\mathbb{R}^n$  of eigenvectors of  $A$ , i.e.

$$A\eta_i = \lambda_i \eta_i \quad i=1, \dots, n, \text{ since } \eta_1, \dots, \eta_n \text{ are a basis,}$$

$$x_0 = \sum_{i=1}^n \alpha_i \eta_i, \quad \alpha_i \in \mathbb{R} \text{ and } \alpha_1 = \eta_1^T x_0 \neq 0$$

$$\begin{aligned} x_k &= A x_{k-1} = A^k x_0 = \sum_{i=1}^n \alpha_i A^k \eta_i = \sum_{i=1}^n \alpha_i \lambda_i^k \eta_i \\ &= \alpha_1 \lambda_1^k \underbrace{\left( \eta_1 + \sum_{i=2}^n \frac{\alpha_i}{\alpha_1} \left( \frac{\lambda_i}{\lambda_1} \right)^k \eta_i \right)}_{z_k} \end{aligned}$$

$$\rightarrow \lim_{k \rightarrow \infty} z_k = \eta_1 \text{ since } |\lambda_2| < |\lambda_1|$$

$$y_k = \frac{x_k}{\|x_k\|} = \frac{\overline{\alpha_1} \lambda_1^k z_k}{\|k z_k\| \lambda_1^k \|z_k\|} \xrightarrow{k \rightarrow \infty} \pm \frac{z_k}{\|z_k\|} \xrightarrow{k \rightarrow \infty} \pm \eta_1$$

for  $k \rightarrow \infty$ .  $\square$

Remarks: •) This method only allows to find the largest eigenvalue & eigenvector.

If I have the eigenvector  $\eta_1$ , know that

$$\begin{aligned} A\eta_1 &= \lambda_1 \eta_1 \\ \Rightarrow \eta_1^T A \eta_1 &= \lambda_1 \eta_1^T \eta_1 \\ \Rightarrow \lambda_1 &= \frac{\eta_1^T A \eta_1}{\eta_1^T \eta_1} \quad \text{Rayleigh coefficient} \end{aligned}$$

•) The speed of convergence depends on the gap between  $|\lambda_1|$  and  $|\lambda_2|$ . If that gap is small, e.g.  $\frac{|\lambda_2|}{|\lambda_1|} = 0.9999 \rightarrow$  slow convergence

If gap is larger, e.g.  $\frac{|\lambda_2|}{|\lambda_1|} = 0.7 \rightarrow$  fast convergence.

Q4, HW3:  $A \in \mathbb{R}^{n \times n}$ ,  $\lambda$  is an eigenvalue of  $A^T A$ ,  $x \in \mathbb{R}^n$ ,  $x \neq 0$   
corresponding eigenvector

(a) Show that  $\|Ax\|_2^2 = \lambda \|x\|_2^2$  and thus  $\lambda \geq 0$ .

$$\underbrace{x^T A^T A x}_{\|Ax\|_2^2} = \underbrace{\lambda x^T x}_{\lambda \|x\|_2^2}$$

(b)  $\|\cdot\|$  norm in  $\mathbb{R}^n$

$$\|A\|_2 \leq \|A^T A\|^{\frac{1}{2}}$$

$$A^T A x = \lambda x \rightarrow \|A^T A x\| = |\lambda| \|x\| = \lambda \|x\|$$

$$\rightarrow \lambda \leq \|A^T A\|$$

$$\|A\|_2 = \max_{1 \leq i \leq n} \sqrt{\lambda_i(A^T A)} \leq \|A^T A\|^{\frac{1}{2}}$$

eigenvalues of  $A^T A$ .

(c)

$$\|A\|_2 \leq \|A^T A\|^{\frac{1}{2}} \leq \|A\|^{\frac{1}{2}} \|A^T\|^{\frac{1}{2}}, \kappa_2(A) \leq (\kappa_1(A) \kappa_\infty(A))^{\frac{1}{2}}$$

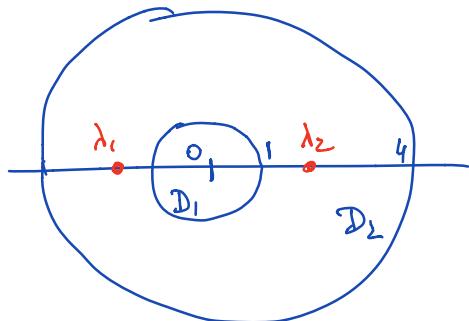
$$= (\|A\|^{\frac{1}{2}}, \|A\|_\infty^{\frac{1}{2}}) \|A\|_2 \|A^{-1}\|_2 \leq (\|A\|, \|A\|, \|A\|_\infty, \|A^{-1}\|_\infty^{\frac{1}{2}})$$

$$\|A^{-1}\|_2 \leq \|A\|^{-\frac{1}{2}}, \|A\|_\infty^{-\frac{1}{2}}$$

Gerschgorin:

$$A = \begin{bmatrix} 0 & 1 \\ 4 & 0 \end{bmatrix}$$

eigenvalues are roots of  
 $\lambda^2 - 4 \rightarrow \lambda_{1,2} = \pm 2$



Can the number 6.125 be represented exactly in a computer?

$$4 + 2 + \frac{1}{8} \text{ in binary: } 110.001$$

$\frac{1}{10}$  cannot be represented: 0.0001001001...

Example when Newton does not converge quadratically  
 for a  $C^2$  function: At solution  $\underline{f'(x^*) \neq 0}$

Example: Solve  $f(x) = x^3 = 0$ ,  $f'(x) = 3x^2$  at

Solution  $f'(x^*) = f'(0) = 0$ .

Assumption:  $\left| \frac{f''(x)}{f'(y)} \right| \leq A \quad \forall x, y \text{ close to solution}$

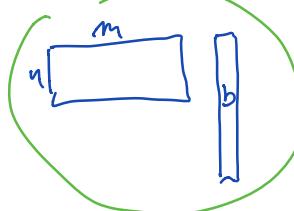
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Least squares problem  $Ax = b$   $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$

$$A = \hat{Q} \hat{R}$$

$$\boxed{\quad} = \boxed{\quad} \boxed{\cancel{\quad}}$$

How many flops does it take to solve  
the least squares problem?

$$\hat{R}^T x = \hat{Q}^T b$$


$n$   
 $m^2$

$m$   
 $n$

$b$

$m$  mult  
 $m-1$  addition }  $\times n$ -dimes

$m(2m-1)$

$$m^2 + 2mn + \dots$$