

## 5. Markov Processes.

A stochastic process in discrete time is just a sequence  $\{X_j : j \geq 0\}$  of random variables with values in some  $(\mathcal{X}, \mathcal{F})$  defined on a probability  $(\Omega, \Sigma, P)$ . It can also be specified by prescribing, in a self consistent manner, the joint distribution of  $\{X_0, X_1, X_2, \dots, X_n\}$  for every  $n$ . A convenient way of doing it is by specifying the the distribution  $p_0(dx_0)$  of  $X_0$  and the conditional distributions

$$p_n(x_0, x_1, \dots, x_{n-1}; dx_n)$$

of  $X_n$  given  $X_0, \dots, X_{n-1}$ .  $\Omega$  can be the product space  $\mathcal{X}^\infty$ , i.e. the space of sequences with values in  $\mathcal{X}$ . There is a canonical  $P$  on the natural  $\sigma$ -field  $\mathcal{F}_\infty$  on  $\Omega$ . There is also the sub- $\sigma$ -fields  $\mathcal{F}_n$  generated by  $x_0, x_1, \dots, x_n$ . The canonical  $P$  will equal  $p_0$  on  $\mathcal{F}_0$  and the conditional distribution of on  $\mathcal{F}_n$  given  $\mathcal{F}_{n-1}$  will be given by  $p_n(x_0, x_1, \dots, x_{n-1}; dx_n)$ . In the special case when  $p_n(x_0, x_1, \dots, x_{n-1}; dx_n) = \pi_n(x_{n-1}, dx_n)$ , for  $n \geq 1$ , depends only on  $x_{n-1}$ , the process is called a Markov Process. Of course when they are just  $p_n(dx_n)$  and do not depend on any  $x_i$  for  $0 \leq i \leq n-1$  we have independent random variables and  $P$  is the product measure. If, in the Markov case,  $\pi_n(\cdot, \cdot)$  is the same  $\pi(\cdot, \cdot)$  for all  $n \geq 1$ , it is called a Markov process with stationary transition probabilities.

A simple example is to take  $\mathcal{X}$  to be a countable set. Then  $p_0$  is just the set of probabilities  $p_0(x) = P[X_0 = x]$ , and

$$P[X_n = y | X_0 = x_0, \dots, X_{n-1} = x_{n-1}] = \pi_n(x, y)$$

are the transition probabilities which in the stationary case is independent of  $n$ . It is natural to consider  $(\Omega, \mathcal{F}_n, \mathcal{F}_\infty, P)$ . There are some natural martingales. For simplicity we limit ourselves to the stationary case.

**Theorem.** For any function  $f$  on  $\mathcal{X}$  let us define

$$(\pi f)(x) = \sum_y \pi(x, y) f(y) = E[f(X_n) | X_{n-1} = x]$$

Then

$$Z_n = f(X_n) - f(X_0) - \sum_{j=0}^{n-1} (\pi f - f)(X_j)$$

is a martingale with respect to  $(\Omega, \mathcal{F}_n, P)$ .

**Proof:** Let us compute  $E[Z_n | \mathcal{F}_{n-1}]$ .

$$\begin{aligned}
 E[Z_n | \mathcal{F}_{n-1}] &= E[f(X_n) | \mathcal{F}_{n-1}] - \sum_{j=0}^{n-1} (\pi f - f)(X_j) \\
 &= (\pi f)(X_{n-1}) - \sum_{j=0}^{n-1} (\pi f - f)(X_j) \\
 &= f(X_{n-1}) - \sum_{j=0}^{n-2} (\pi f - f)(X_j) \\
 &= Z_{n-1}
 \end{aligned}$$

**Remark.** If we replace the definition  $(\pi f)(x) = \sum_y \pi(x, y) f(y)$  with

$$(\pi f)(x) = \int f(y) \pi(x, dy)$$

then the theorem is true for Markov processes on any state space. For simplicity we will assume that we have a countable state space.

Martingales are a useful tool in studying Markov Processes. Let us look at some examples.

1. Let  $A \subset \mathcal{X}$ . Define

$$\tau_A = \inf\{j : X_j \in A\}$$

is the first hitting time of  $A$ . It is possible that  $X_j$  never hits  $A$  in which case we take  $\tau_A = \infty$ . We wish to calculate for  $\lambda > 0$ ,

$$(5.1) \quad \phi_\lambda(x) = E[e^{-\lambda \tau_A} | X_0 = x]$$

Then if  $x \in A$  then  $\phi_\lambda(x) = 1$ . Moreover for  $x \notin A$  it is easy to see that

$$\phi_\lambda(x) = e^{-\lambda} \sum_y \pi(x, y) \phi_\lambda(y)$$

Clearly  $0 \leq \phi_\lambda(x) \leq 1$ . We will show that the only bounded solution of

$$(5.2) \quad F(x) = e^{-\lambda} \sum_y \pi(x, y) F(y)$$

for  $x \notin A$  with  $F(x) = 1$  for  $x \in A$  is given by (5.1). Let  $F(x)$  be a solution of (5.2). Define

$$Z_n = e^{-\lambda n} F(X_n)$$

Then with  $\Sigma_n = \sigma\{X_0, X_1, \dots, X_n\}$ ,

$$\begin{aligned} E[Z_{n+1} | \Sigma_n] &= e^{-\lambda(n+1)} E[F(X_{n+1}) | \Sigma_n] \\ &= e^{-\lambda(n+1)} \sum_y \pi(X_n, y) F(y) \\ &= e^{-\lambda n} F(X_n) \\ &= Z_n \end{aligned}$$

provided  $X_n \notin A$ . One can rewrite this as

$$E[Z_{n+1} - Z_n | \sigma_n] = \begin{cases} 0 & \text{if } X_n \notin A \\ e^{-\lambda n} G(X_n) & \text{if } X_n \in A. \end{cases}$$

with

$$G(x) = e^{-\lambda} \sum_y \pi(X_n, y) F(y) - F(x)$$

Therefore

$$Z_n - Z_0 - \sum_{j=0}^{n-1} e^{-\lambda j} G(X_j) \mathbf{1}_A(X_j)$$

is a martingale. Let  $\tau_A$  is a stopping time and for  $n \leq \tau_A$   $X_n \notin A$  and  $G(X_n) = 0$ . Therefore  $\{Z_n\}$  is bounded uniformly until  $\tau_A$  even if  $\tau_A$  itself can be large. Doob's stopping theorem applies and

$$E[e^{-\lambda \tau_A}] = E[e^{-\lambda \tau_A} F(X_{\tau_A})] = E[Z_\tau] = E[Z_0] = F(x)$$

**Example.** Consider the random walk on  $Z$  where  $\pi(x, x \pm 1) = \frac{1}{2}$ . If one starts from 0, and  $\tau$  is the first time  $\pm k$  is reached calculate  $E[e^{-\lambda \tau}]$ . Solve the equation

$$F(x) = e^{-\lambda} \left[ \frac{1}{2} F(x-1) + \frac{1}{2} F(x+1) \right]$$

for  $|x| \leq k - 1$ , with  $F(x) = 1$  for  $|x| \geq k$ . One can isolate  $[-k, k]$ . Need to solve

$$F(x - 1) + F(x + 1) - 2e^\lambda F(x) = 0$$

with  $F(\pm k) = 1$ . Solve the quadratic

$$\rho^2 - 2e^\lambda \rho + 1 = 0$$

with roots

$$\rho_{\pm} = e^\lambda \pm \sqrt{e^{2\lambda} - 1} = e^{\pm\theta}$$

where  $\theta = \log[e^\lambda + \sqrt{e^{2\lambda} - 1}]$ . The solution is seen to be

$$F(x) = \frac{e^{\theta x} + e^{-\theta x}}{e^{\theta k} + e^{-\theta k}}$$

and

$$F(0) = [\cosh(\theta k)]^{-1}$$

**Exercise.** Start from  $x > 0$ . Show that sooner or later 0 is reached. Calculate  $E[e^{-\lambda \tau}]$  where  $\tau$  is the first time 0 is reached.

**Exercise.** What happens when

$$p = \pi(x, x - 1) > \frac{1}{2} > \pi(x, x + 1) = q = 1 - p$$

and when

$$p = \pi(x, x - 1) < \frac{1}{2} < \pi(x, x + 1) = q = 1 - p$$

**Example.** A game is being played where the probability is  $\frac{1}{2}$  for each of two players to win any one round. It is agreed that the first person to win  $k$  rounds will be the winner. They put equal amounts to make a kitty for the winner to take. Unfortunately the game is interrupted before either player can win  $k$  rounds. It stops when player A needs to win  $a$  more rounds, and player B needs  $b$  rounds.  $1 \leq a \leq k$ ,  $1 \leq b \leq k$ . What is the "fair" way to divide the kitty between the two players?

Let  $u(a, b)$  be the proportion of the kitty that player A should get in a fair division when he needs  $a$  rounds and player B needs  $b$  rounds. Since the

game is fair neither player can expect to gain or lose by playing an extra game.

$$u(a, b) = \frac{1}{2}u(a - 1, b) + \frac{1}{2}u(a, b - 1)$$

$u(0, b) = 1$  if  $b > 0$  and  $u(a, 0) = 0$  if  $a > 0$ . Solution is

$$u(a, b) = \frac{1}{2^{a+b-1}} \sum_{a+b-1 \geq r \geq a} \binom{a+b-1}{r}$$

You can verify that this is a solution. Can you show directly ?